Performance Evaluation of Zigbee Protocol for Habitat Monitoring.

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ABSTRACT:-

Habitat and environmental monitoring are driving applications for wireless sensor network. In this paper we present a performance evaluation of Zigbee which is IEEE 802.15.4 standard, focusing on energy required for sending and receiving uplink packets and hence battery lifetime. We analyzed the performance based on these two metrics. The analytical and simulation results are perfectly matched.

KeyWords: WSN, Habitat, Zigbee, MAC, CSMA.

I. INTRODUCTION: Wireless adhoc network is a generic term grouping different networks, which are self organizing, meaning that there is neither a centralized administration nor a fixed network infrastructure, and communication links are wireless. A wireless sensor network is an adhoc network consisting of spatially distributed autonomous sensor nodes which are equipped with radio transceiver, a microcontroller, an energy source (usually a battery) and sensor to cooperatively monitor or environmental conditions.

In wireless sensor networks, maximizing battery lifetime is very important design criterion, because in many applications charging or recharging battery after deployment is normally not economical or not feasible.

Current trend in computer communication is gearing towards embedded ubiquitous computing. Advances in wireless communications and microelectronics enabled the development of low cost sensors which are constantly being adapted for pervasive communication paradigm. One of the main objectives of the MAC protocol for WSN is to minimize power consumption while providing reliable, low data rate transmission.

Unlike other wireless network, it is generally very hard to charge/replace the battery, which gives the way to the primary objective of maximizing network lifetime, leaving the other performance metrics as secondary objective. Since the communication of sensor nodes will be more energy consuming than their computations, it is primary concern that communications is minimized whole achieving the desired network operation.

Habitat and environmental monitoring represent a class of sensor network applications with

enormous potential benefits for scientific communities and society as whole. applications share a common structure, where field of sensors are tasked to take periodic reading and report the result and derived values to a central repository. In the context of habitat monitoring, wireless sensor networks offers significant advantages. Individual devices can be made significantly numerous to take measurements at many locations of interest and mitigate errors arising from the interpolation and extrapolation from coarser-grained samples. They can be sufficiently small to be co-located with phenomena of interest without altering the parameters to be measured. And they can be unobtrusively embedded in the environment without creating conspicuous landmarks that

This paper presents a comparison of analytical and simulation model of Zigbee, which is low power MAC. This paper describes comparison results of energy consumption required for sending and receiving packets as well as lifetime of battery.

change the behaviors of its inhabitants.

II. RELATED WORK: The smooth operation of any wireless network depends, to a large extent, on the effectiveness of the low-level Medium Access Control (MAC) layer responsible for sharing the ether. A MAC protocol determines the next node to access the medium, tries to ensure that no two nodes are interfering with each other's transmissions, and deals with the situation when they do.

TDMA-based protocols, in contrast to contention-based protocols, are very effective at avoiding collisions and have a built-in duty cycle mitigating idle listening. They require, however, some authority (e.g., a dedicated access point) to orchestrate activities within a cell. This complicates their deployment in multi-hop ad-hoc (sensor) networks where nodes are equal and have limited resources.

The second approach mitigating idle listening is the S-MAC protocol, a true MAC protocol, which also addresses the overheads caused by collisions, overhearing, and protocol overhead. The basic idea of this contention-based protocol is that time is divided into relatively large

frames. Every frame has two parts: an active part and a sleeping part. During the sleeping part, a node turns off its radio to preserve energy.

During the active part, it can communicate with its neighbors and send any messages queued during the sleeping part on idle listening, is reduced. The exact savings are under control of the application. Consequently, it results in increase in latency, and reduction in throughput. S-MAC needs some synchronization between nodes, but that is not as critical as in TDMA-based protocols. The S-MAC protocol uses the RTS/CTS/DATA/ACK signaling scheme from 802.11 to reduce the number of collisions caused by the hidden-node problem. S-MAC includes message passing support to reduce protocol overhead when streaming a sequence of message fragments.

An energy-saving protocol for sensor networks considered here is the T-MAC protocol, which automatically adapts the duty cycle to the network traffic. As with S-MAC, nodes form a virtual cluster to synchronize themselves on the beginning of a frame. But instead of using a fixed-length active period, T-MAC uses a time-out mechanism to dynamically determine the end of the active period. If receiver does not receive any message during time out interval, it goes to sleep and if it receives such message, the time start starts afresh after reception of message. The basic T-MAC scheme suffers from so called early sleep problem which can reduce throughput.

B-MAC, developed at the University of California at Berkeley, is a CSMA-based technique that utilizes low power listening and an extended preamble to achieve low power communication. Nodes have an awake and a sleep period, and each node can have an independent schedule. If a node wishes to transmit, it precedes the data packet with a preamble that is slightly longer than the sleep period of the receiver. During the 'awake' period, a node samples the medium and if a preamble is detected it remains awake to receive the data. With the extended preamble, a sender is assured that at some point during the preamble the receiver will wake up, detect the preamble, and remain awake in order to receive the data. B-MAC also provides an interface by which the application can adjust the sleep schedule to adapt to changing traffic loads. The method of adaptation is left to the application developer. It is shown that B-MAC surpasses existing protocols in terms of throughput, latency, and for most cases energy consumption. While B-MAC performs quite well, it suffers from the overhearing problem, and the long preamble dominates the energy usage.

WiseMAC, which is based on ALOHA, also uses preamble sampling to achieve low power communications in infrastructure sensor networks. WiseMAC uses a similar technique to B-MAC, but the sender learns the schedules of the receiver awake periods, and schedules its transmission so as to

reduce the length of the extended preamble. In addition, for low traffic loads where the preamble is longer than the data frame, WiseMAC repeats the data frame in place of the extended preamble.

Receivers process this data frame and if the node is not the intended recipient it returns to sleep. If the node is the recipient, it remains awake until the end of the transmission and sends an acknowledgement. While WiseMAC solves many of the problems associated with low power communications, it does not provide a mechanism by which nodes can adapt to changing traffic patterns.

III. IDENTIFICATION OF PROBLEM AND ORGANIZATION:

As WSN are useful in many monitoring applications, in this paper we concentrating one of the applications of WSN, which is Habitat monitoring. For Habitat monitoring following as the core set of requirements.

- **1. Energy efficiency**: The sensor nodes should not need refreshing of batteries for at least 6 months.
- **2. Remote querying and reconfigurability**: It should be possible to query the data in several formats and to reconfigure the monitoring parameters via the Internet.
- 3. Ease of deployment: The client should be able deploy the system without any need for special configuration.
- 4. Reliability: The provided data should be realtime and high-fidelity. The system should be available 99% of the time and should recover from a crash quickly.

As very prior important parameter for Habitat monitoring is energy efficiency, from the above survey of energy efficient MAC, we have decided to implement ZigBee MAC protocols for Habitat Monitoring application and compare results with respect to energy saving and battery lifetime.

The organization of our work is like this: Part A) in section IV represents analytical model for Zigbee part B) of section IV represents simulation model for Zigbee computation details.

The definition and default values of parameters used for computation are listed in table (I). For simplicity computation is carried out for star topology and for uplink transmission. Poisson process traffic model is used. In this probability

density function is given as, $P(t) = \lambda e^{-tt}$ and cumulative distribution function is, $f(t) = 1 - e^{-tt}$.

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A) Zigbee analytical computation details:
$$P_{sleep} = P_{r_fsleep} + P_{\mu c_sleep} + P_{other} \qquad (1) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (2) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (2) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (2) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (3) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{bcn} / T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{cycle}) + P_{sleep} \qquad (4) \\ P_{no_traffic} = P_m * (T_{s_rf} + T_{cycle})$$

$$E_s = P_{slave} * (CAP + T_{_t} + T_{_data} + T_{_ack})$$

= Expected energy to send a packet from slave to master. (3)

$$\begin{split} E_r &= P_m^* \left(T_{bcn_extra} + T_{_data} + T_{_ack} + T_{_t} \right) + P_{sleep} \\ &= Expected \ energy \ to \ receive \ packet \ from \end{split}$$

slave. (4)

 $P_{oh} = P_{slave} * (T_{bcn_extra})$ = Power consumption for overhearer. (5) $E_{s_avg} = E_s / N + \hat{P}_{oh} * (N-1) / N$ = Average sending power. (6) $E_{r_avg} = E_r + P_{oh} * (N-1) / N$ = Average receiving power. (7) $P_{\lambda_Es}\!=E_{s_avg}*\lambda*N+P_{no_traffic}$ (8) $P_{\lambda \text{ Es days}} = (108/864 * P_{\lambda \text{ Es}})$ (9) $P_{\lambda_{\text{Er}}} = (E_{r_{\text{avg}}} * \lambda) + P_{\text{no_traffic}}$ (10) $P_{\lambda \text{ Er days}} = (108/864 * P_{\lambda \text{ Er}})$ (11)

B)Simulation details:

The definition and default values of parameters used for computation are listed in table (I).

For simplicity computation is carried out for star topology and for uplink transmission. Poisson process traffic model is used. In this probability density function is given as, $P(t) = \lambda e^{-2t}$ and cumulative distribution function is, $f(t) = 1 - e^{-2t}$.

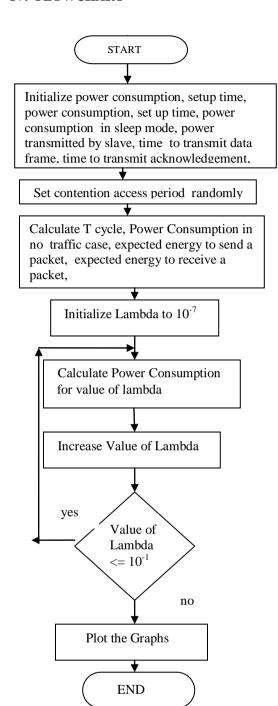
Free-space propagation model is used to detect signal strength. CSMA technique is used to access the medium. Backoff counter is set to maximum value and when it hits to zero medium is accessed.

Table I: Parameters of the WSN.

Parameter	Description	Value
U_{system}	Supply voltage of the System	3 V
B _{capacity}	Capacity of the battery	1000
Deapacity	Capacity of the battery	mAh
fuc	Working frequency of the	9 MHz
$J\mu C$	microcontroller	9 WILL
D	Power consumption of the RF	27 μW
$P_{rf\ sleep}$		27 μW
	in sleep mode	00 117
$P_{\mu c \; sleep}$	Power consumption of the	90 μW
	microcontroller in sleep mode	27 11
P_{others}	Power consumption of other	27 μW
	components: memory, RTC,	
	etc.	
$P_{rf rx}$	Power consumption of the RF	27 mW
	in RX mode	
P_{rftx}	Power consumption of the RF	39.9
	in TX mode	mW
$P_{\mu c}$	Power consumption of the	32.7
	microcontroller in working	mW
	mode	
$T_{to \ rx}$	Setup time of the RF RX mode	2.2 ms
$T_{to tx}$	Setup time of the RF TX mode	1.1 ms
T_t	Transition time of the RF	T _{to rx}
	between RX and TX mode	
$T_{s r s s i}$	Setup time of the RF RSSI	0.4 ms
$T_{s rf}$	Setup time of the RF from	T _{to rx}
3.9	sleep to awake	10 11
$T_{s \mu c}$	Setup time of the	2.6 ms
3 με	microcontroller from sleep to	
	awake	
$T_{samplerssi}$	Time for sampling the channel	0.4 ms
- sampierssi	if only RSSI is used	0111111
R_{data}	Transmission rate for DATA	100
raaia	frame	bit/s
R_{wus}	Transmission rate for wake-	70 bit/s
It _{wus}	up-signal	70 010 8
I	Length of ACK frame	96 bit
L_{ack}		
L_{swuf}	Length of a short wake-up-	326 bit
	frame	576 0
Lip max	Maximum length of IP packet	576 × 8
<u> </u>		bit
$L_{data\ max}$	Maximum length of DATA	7200 bit

	frame	
L_{bcn}	Length of an IEEE 802.15.4	276 bit
	MAC beacon frame	
L _{bcn extra}	Length of an IEEE 802.15.4	96 bit
-ben exira	frame indication in beacon	, , , , , ,
L_{cmd}	Length of an IEEE 802.15.4	228 bit
	MAC data request command	
	frame	
$L_{to addr}$	Length of a frame until the	192 bit
	destination address	
T_{ack}	Time to transmit an ACK	0.96 ms
	frame	
T_{swuf}	Time to transmit an SWUF	4.66 ms
_	frame	
T_{data}	Time to transmit a DATA	-
	frame	
$T_{data\ max}$	Time to transmit a DATA	72 ms
	frame with maximal length	
T_{hcn}	Time to transmit an IEEE	2.76 ms
	802.15.4 beacon frame	
T _{bcn extra}	Time to transmit an IEEE	0.96 ms
	802.15.4 frame indication in	
	beacon	
T_{cmd}	Time to transmit an IEEE	2.28 ms
Cmu	802.15.4 data request	
	command frame	
$T_{to \ addr}$	Time to transmit a frame until	1.92 ms
- 10 daar	the destination address	2., 2 1110
N	Number of slaves	_
		l

IV. FLOWCHART



V.RESULTS:

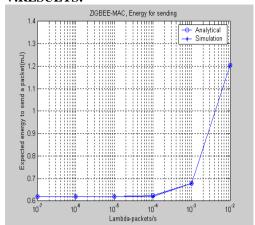


Fig.1-Energy required for sending packets/sec

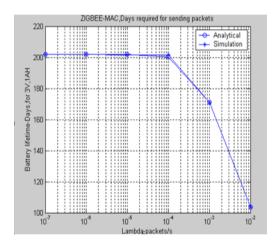


Fig.3 –Battery Life-time required for sending packets/sec.

VI. CONCLUSION:

We have presented comparison of analytical and simulation results of Zigbee. From the results it is seen that ZigBee MAC is suitable for low traffic. The results indicate that the analytical and simulation model agree. Hence the simulation model is validated. The work on this is in progress, it will be interesting to modify this simulation model to improve the energy efficiency of the WSN.

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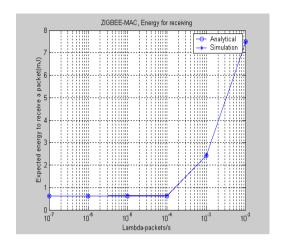


Fig.2-Energy required for receiving packets/sec

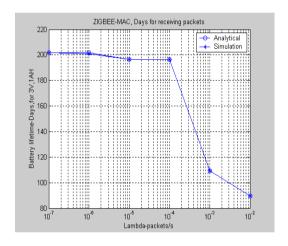


Fig.4- Battery Life-time required for receiving packets/sec.

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