Secure AODV Simulation under Black hole Attack in MANET

Vidhya Patel
M. Tech Scholar
Computer Science Engineering Department
IET, Alwar, Rajasthan

Deepak Chaudhary
Assistant Professor
Computer Science Engineering Department
IET, Alwar, Rajasthan

ABSTRACT
This paper dissects the blackhole assault which is one of the conceivable assaults in impromptu systems. In a blackhole assault, a malignant hub mimics a goal hub by sending a ridiculed course answer bundle to a source hub that starts a course revelation. By doing this, the noxious hub can deny the activity from the source hub. In this paper, we mimic the Ad hoc on Demand Vector Routing Protocol (AODV) under blackhole assault by thinking about various execution metric. The recreation results demonstrate the viability of blackhole assault on AODV convention.

Keywords
AODV, Blackhole attack, Security, MANET

1. INTRODUCTION
Versatile impromptu system (MANET) is an accumulation of portable hosts without the required intercession of any current framework or incorporated passageway, for example, a base station. The uses of MANET extend from an erratic gathering system, crisis activities, for example, fiasco recuperation to military applications because of their simple arrangement. In any case, because of their inborn qualities of dynamic topology and absence of unified administration security, MANET is helpless against different sorts of assaults. These incorporate inactive listening stealthily, dynamic meddling, pantomime, and foreswearing of-benefit Blackhole assault is one of numerous conceivable assaults in MANET. In this assault, a pernicious hub sends a fashioned Route REPLY (RREP) bundle to a source hub that starts the course disclosure with a specific end goal to claim to be a goal hub. By looking at the goal arrangement number contained in RREP bundles when a source hub got different RREP, it judges the best one as the latest steering data and chooses the course contained in that RREP parcel. In the event that the grouping numbers are equivalent it chooses the course with the littlest jump tally. On the off chance that the assailant parodied the character to be the goal hub and sends RREP with goal grouping number higher than the genuine goal hub to the source hub, the information movement will stream toward the aggressor. In this way, source and goal hubs ended up unfit to speak with each other. In [14], the creators

Investigated the effect of blackhole attack when movement velocity and a number connection toward the victim node are changed, and proposed the detection technique at the destination node. In Section 2 of this paper, we discuss Related work. In Section 3, we summarize the basic operation of AODV (Ad hoc On-Demand distance Vector Routing) protocol on which we base our work. In Section 4, we describe the effect of blackhole attack on AODV. Section 5 presents the simulation of AODV under blackhole attack. Section 6 discusses the simulation result based on simulation experiments. Finally, Section 7 presents conclusion.

2. RELATED WORK
There without a doubt have been various endeavors distributed in the writing that go for countering the Black opening assaults. We study them in the accompanying. In [5], the creators talk about a convention that requires the middle of the road hubs to send RREP message alongside the following jump data. At the point when the source hub gets this data, it sends a RREQ to the following jump to confirm that the objective hub (i.e. the hub that simply sent back the RREP bundle) undoubtedly has a course to the transitional hub and to the goal. At the point when the following jump gets a Further Request, it sends a Further Reply which incorporates the check result to the source hub. In view of data in Further Reply, the source hub judges the legitimacy of the course. In this convention, the RREP control bundle is changed to contain the data about next jump. In the wake of accepting RREP, the source hub will again send RREQ to the hub indicated as next bounce in the got RREP. Clearly, this builds the steering overhead and end-to-end delay. What's more, the middle of the road hub needs to send RREP message twice for a solitary course ask. In [6], the creators portray a convention in which the source hub checks the credibility of a hub that starts RREP by discovering in excess of one course to the goal. At the point when source hub gets RREPs, if courses to goal shared jumps, source hub can perceive a protected course to goal. Sanjay Ramaswamy, et al [7] proposed a technique for recognizing various dark opening hubs. They are first to propose answer for helpful dark gap assault. They somewhat altered AODV convention by presenting information directing data table (DRI) and cross checking. Each section of the hub is kept up by the table. They depend on the solid hubs to exchange the bundles. In [8] proposed an answer with the improvement of the AODV convention which stays away from different dark openings in the gathering. A method is provide for recognize different dark gaps collaborating with each other and find the sheltered course by maintaining a strategic distance from the assaults. It was expected in the arrangement that hubs are now confirmed and consequently can take an interest in the correspondence. It utilizes Fidelity table where each hub that is taking an interest is given a loyalty level that will give dependability to that hub. Any node having 0 values is considered as malicious node and is eliminated.

In [9] proposed the solution which discovers the secure route between source and destination by identifying and isolating cooperative black hole nodes. This solution adds on some changes in the solution proposed by the S.Ramaswamy to improve the accuracy. This algorithm uses a methodology to identify multiple black hole nodes working collaboratively as a group to initiate cooperative black hole attacks.

3. THEORETICAL BACKGROUND OF AODV

AODV is a reactive routing protocol; that do not lie on active paths neither maintain any routing information nor participate in any periodic routing table exchanges. Further, the nodes do not have to discover and maintain a route to another node until the two needs to communicate, unless former node is offering its services as an intermediate forwarding station to maintain connectivity between other nodes [2]. AODV has borrowed the concept of destination sequence number from DSDV [5], to maintain the most recent routing information between odes.. Whenever a source node needs to communicate with another node for which it has no routing information, Route Discovery process is initiated by broadcasting a Route Request (RREQ) packet to its neighbors. Each neighboring node either responds the RREQ by sending a Route Reply (RREP) back to the source node or rebroadcasts the RREQ to its own neighbors after increasing the hop_count field. If a node cannot respond by RREP, it keeps track of the routing information in order to implement the reverse path setup or forward path setup. The destination sequence number specifies the freshness of a route to the destination before it can be accepted by the source node. Eventually, a RREQ will arrive to node that possesses a fresh route to the destination. If the intermediate node has a route entry for the desired destination, it determines whether the route is fresh by comparing the destination sequence number in its route table entry with the destination sequence number in the RREQ received. The intermediate node can use its recorded route to respond to the RREQ by a RREP packet, only if, the RREQ's sequence number for the destination is greater than the recorded by the intermediate node.

Instead, the intermediate node rebroadcasts the RREQ packet. If a node receives more than one RREPs, it updates its routing information and propagates the RREP only if RREP contains either a greater destination sequence number than the previous RREP, or same destination sequence number with a smaller hop count. It restrains all other RREPs it receives. The source node starts the data transmission as soon as it receives the first RREP, and then later updates its routing information of better route to the destination node. Each route table entry contains the following information:

• Destination node

• Next hop

• Number of hops

• Destination sequence number

• Active neighbors for the route

• Expiration timer for the route table entry

The route discovery process is reinitiated to establish a new route to the destination node, if the source node moves in an active session. As the link is broken and node receives a notification, and Route Error (RERR) control packet is being sent to all the nodes that uses this broken link for further communication. And then, the source node restarts the discovery process.

As the routing protocols typically assume that all nodes are cooperative in the coordination process, malicious attackers can easily disrupt network operations by violating protocol specification. This paper discusses about blackhole attack and provides routing security in AODV by purging the threat of blackhole attacks

4. DESCRIPTION OF BLACKHOLE ATTACK

MANETs are vulnerable to various attacks. General attack types are the threats against Physical, MAC, and network layer which are the most important layers that function for the routing mechanism of the ad hoc network. Attacks in the network layer have generally two purposes: not forwarding the packets or adding and changing some parameters of routing messages; such as sequence number and hop count. A basic attack that an adversary can execute is to stop forwarding the data packets. As a result, when the adversary is selected as a route, it denies the communication to take place. In blackhole attack, the malicious node waits for the neighbors to initiate a RREQ packet. As the node receives the RREQ packet, it will immediately send a false RREP packet with a modified higher sequence number. So, that the source node assumes that node is having the fresh route towards the destination. The source node ignores the RREP packet received from other nodes and begins to send the data packets over malicious node. A malicious node takes all the routes towards itself. It does not allow forwarding any packet anywhere. This attack is called a blackhole as it swallows all objects; data packets. Fig. 1 Blackhole attacks in MANETs In figure 1, source node S wants to send data packets to a destination node D in the network. Node M is a malicious node which acts as a blackhole. The attacker replies with false reply RREP having higher modified sequence number. So, data communication initiates from S towards M instead of D.

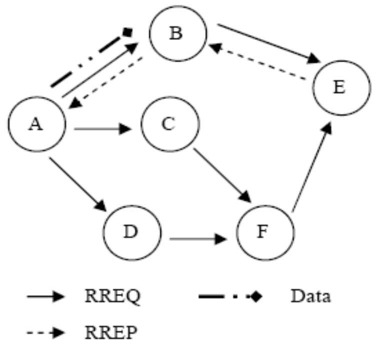


Fig: 1 Blackhole Attack

5. SIMULATION OF AODV UNDER BLACKHOLE ATTACK

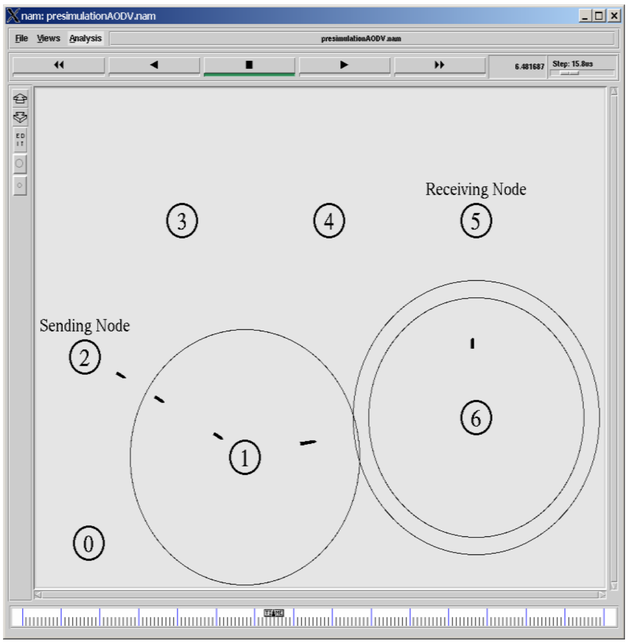
For simulation, we set the parameter as shown in Table 1. Random Waypoint Model (RWP) [1] is used as the mobility model of each node. In this model, each node chooses a random destination within the simulation area and a node moves to this destination with a random velocity. The simulation is done using Network Simulator 2 to analyze the performance of the network by varying the nodes mobility. The metrics used to evaluate the performance are given below.

Packet Delivery Ratio: The ratio between the number of packets originated by the "application layer" CBR sources and the number of packets received by the CBR sink at the final destination.

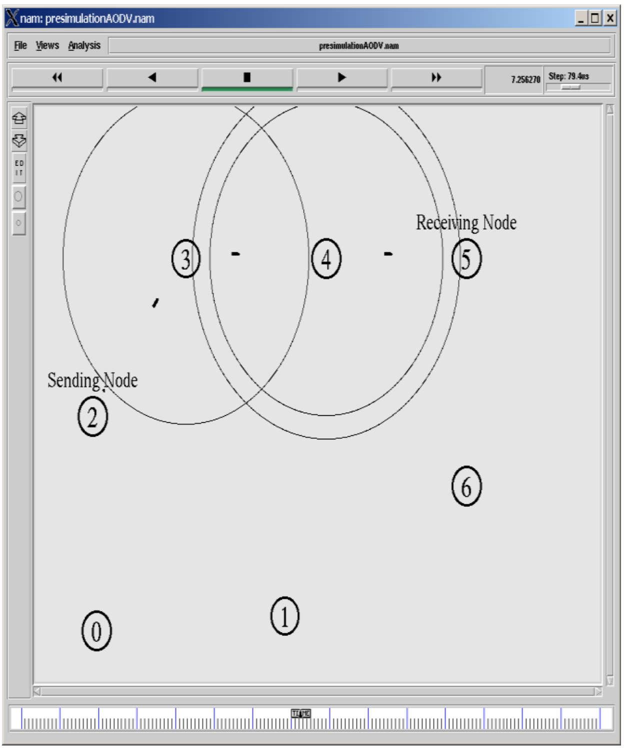
Average End-to-End Delay: This is the average delay between the sending of the data packet by the CBR source and its receipt at the corresponding CBR receiver. This includes all the delays caused during route acquisition, buffering and processing at intermediate nodes, retransmission delays at the MAC layer, etc. It is measured in milliseconds.

Table: 1 Simulation Parameters

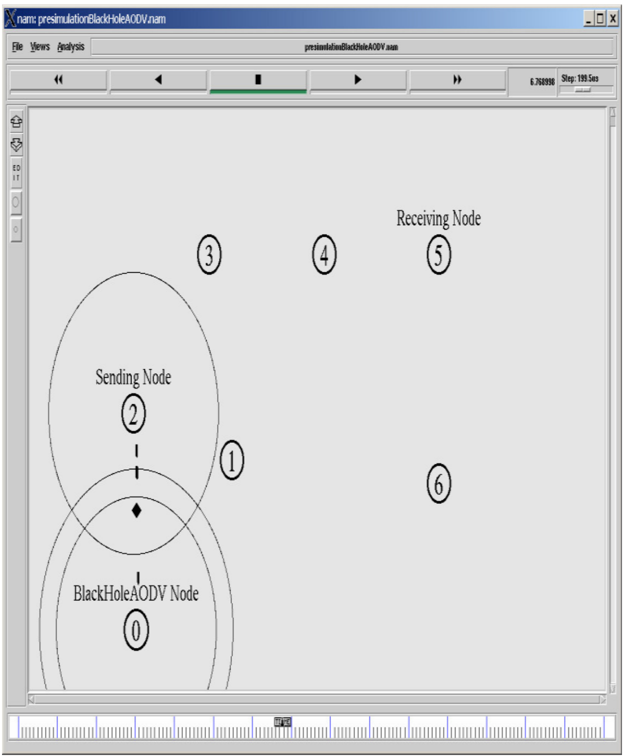
Simulator	Ns-2(version 2.32)
Simulation Time	500 (s)
Number of Mobile Nodes	10, 15,20,25,30
Topology	750 * 750 (m)
Routing Protocol	AODV
Traffic	Constant Bit Rate (CBR)
Pause Time	10 (m/s)
Max Speed	20 (m/s)



Data Flow between Node 2 and 5 via Node 1 & 6



Data Flow between Node 2 & 5 via Node 3 and 4



Node 0 Absorb the Connection Node 2 to Node 5

Here, we assume that the blackhole attack take place after the attacking node received RREQ for the destination node that it is going to impersonate. We also assume that the communication started from a source node to a destination node and the node numbers of the source node, destination node and attacking node are 0, 1 and 9, respectively, as shown in Figure 5 (for 10 nodes). We have carried out the simulation by considering the different number of nodes 10, 15, 20, 25, and 30.

First, we investigate the packet delivery ratio of packet from source node 0 to destination node 1 in case there are connections from other nodes to the destination node. For the experiment, in Figure 2, nodes which are selected randomly from 2 to 8 (for 10 nodes), 2 to 18 (for 20 nodes) etc. (except the source node, destination node, and attacking node) generate traffic toward the destination node. Here, we perform experiment by changing the number of nodes generating the traffic from one to nine.

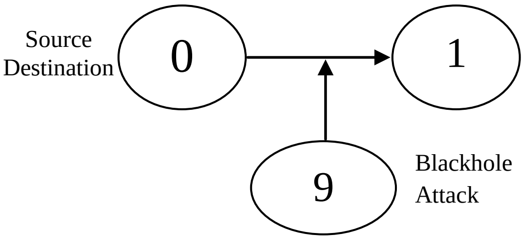


Fig: 2 Node Descriptions

From Figure 3, we can see that when the number of connection is 1, the more Dst Seq is increased in blackhole attack the more packet delivery ratio drops. However, when the number of connections increases, the packet ratio increases even when blackhole attack took place. This is because the destination node's Dst Seq tends to be higher than the attacker's Dst Seq, since attacker set the Dst Seq based on the Dst Seq contained in RREQ coming from the source node. We can see that the more the attacker increases the Dst Seq, the lower the packet delivery rate is.

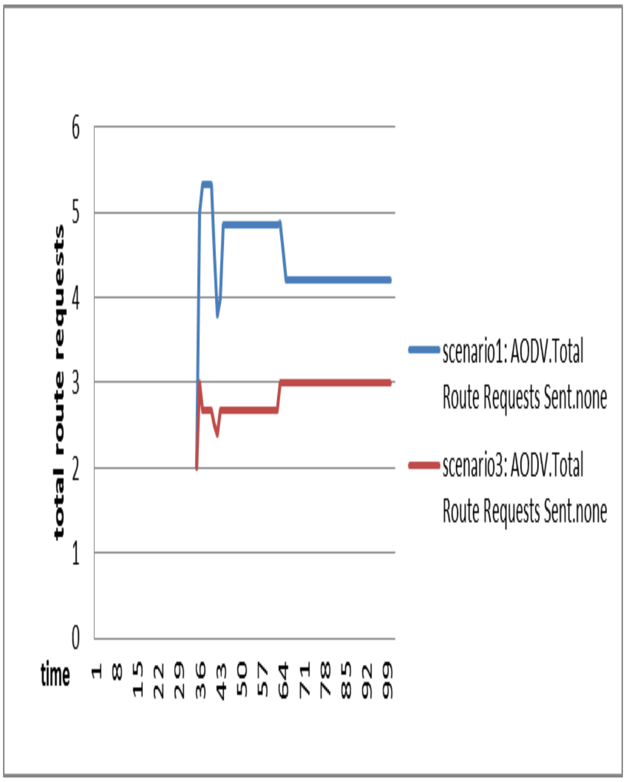


Fig: 3 Packet Delivery Ratios

6. SIMULATION RESULT AND DISCUSSION

Figure-3 shows the packet delivery ratio of normal AODV protocol and in the presence of blackhole attack. In AODV the packet delivery ratio is reduced to 80%. From this figure 3 it is clear that when the malicious node is present in the network, it reduce the packet delivery to destination. From the figure-4 it can be observed that, when blackhole attack initiates in network, there is nearly 21% increase in the average end-to-end delay.

7. CONCLUSION

Blackhole attack is one of the most important security problems in MANET. It is an attack that a malicious node impersonates a destination node by sending forged RREP to a source node that initiates route discovery, and consequently deprives data traffic from the source node. In this paper, we have analyzed the effect of blackhole attack on AODV protocol. The result shows significant degradation in performance of ad hoc on demand vector routing protocol (AODV) under blackhole attack.

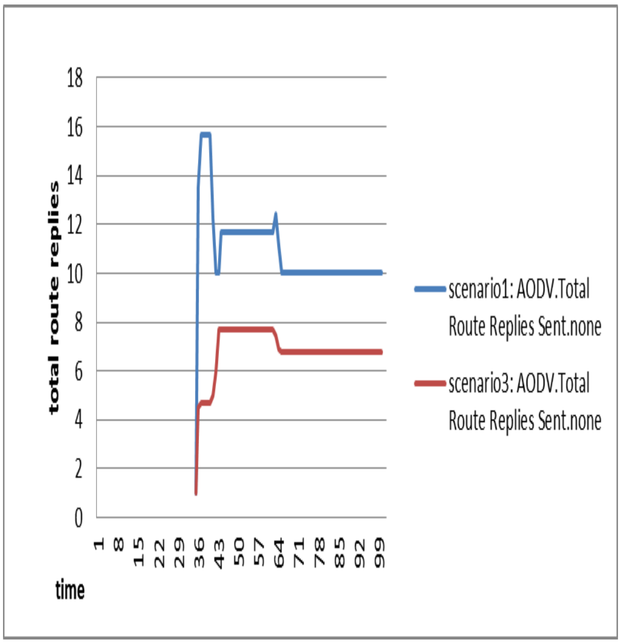


Fig: 4 Average End-to-End Delays

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