# The $L(2,1)$-Labeling on $\gamma$-Product of Graphs and Improved Bound on the $L(2,1)$-Number of $\gamma$-Product of Graphs 

Anuj Kumar<br>Department of Mathematics and Statistics, Gurukula Kangri Vishwavidyalaya Haridwar (U.K), India

P. Pradhan<br>Department of Mathematics and Statistics, Gurukula Kangri Vishwavidyalaya Haridwar (U.K), India


#### Abstract

The concept of $\mathrm{L}(2,1)$-labeling in graph come into existence with the solution of frequency assignment problem. In fact, in this problem a frequency in the form of nonnegative integers is to assign to each radio or TV transmitters located at various places such that communication does not interfere. This frequency assignment problem can be modeled with vertex labeling of graphs. An $\mathrm{L}(2,1)$-labeling (or distance two labeling) of a graph $G$ is a function $f$ from the vertex set $V(G)$ to the set of all nonnegative integers such that $|f(u)-f(v)| \geq 2$ if $d(u, v)=1$ and $|f(u)-f(v)| \geq 1$ if $d(u, v)=2$, where $d(u, v)$ denotes the distance between $u$ and $v$ in $G$. The $\mathrm{L}(2,1)$-labeling number $\lambda(\mathrm{G})$ of G is the smallest number k such that $G$ has an $L(2,1)$-labeling with $\max \{\mathrm{f}(\mathrm{v}): \mathrm{v} \in \mathrm{V}(\mathrm{G})\}=\mathrm{k}$. In this paper, upper bound for the $\mathrm{L}(2,1)$-labeling number for the $\gamma$-product of two graphs has been obtained in terms of the maximum degrees of the graphs involved and improved this bound by using a dramatically new approach on the analysis of the adjacency matrices of the graphs. By the new approach, we have achieved more accurate result with significant improvement of this bound.


## Keywords

Channel assignment, $L(2,1)$-labeling, $L(2,1)$-labeling number, Graph $\gamma$-product, Adjacency matrix of graphs

## 1. INTRODUCTION

The frequency assignment problem asks for assigning frequencies to transmitters in a broadcasting network with the aim of avoiding undesired interference. Hale [20] was first person who formulated this problem as a graph vertex coloring problem. By Roberts [7], In order to avoid interference, any two "close" transmitters must receive different channels and any two "very close" transmitters must receive channels that are at least two channels apart. To translate the problem into the language of graph theory, the transmitters are represented by the vertices of a graph; two vertices are "very close" if they are adjacent and "close" if they are of distance two in the graph. Based on this problem, Griggs and Yeh [10] considered an $L(2,1)$ labeling on a simple graph. An $L(2,1)$-labeling (or distance two labeling) of a graph $G$ is a function $f$ from the vertex set $V(G)$ to the set of all nonnegative integers such that $|f(u)-f(v)| \geq 2$ if $d(u, v)=1$ and $\mid f(u)-$ $f(v) \mid \geq 1$ if $d(u, v)=2$, where $d(u, v)$ denotes the distance between $u$ and $v$ in $G$. A $k-L(2,1)$-labeling is an $L(2,1)$-labeling such that no label is greater than $k$. The $L(2,1)$-labeling number of $G$, denoted by $\lambda(G)$ or $\lambda$, is
the smallest number $k$ such that $G$ has a $k-$ $L(2,1)$ labeling. The $L(2,1)$-labeling has been extensively studied in recent past by many researchers [see $1,4,8,9,11,12,19]$. The common trend in most of the research paper is either to determine the value of $L(2,1)$-labeling number or to suggest bounds for particular classes of graphs.

Griggs and Yeh [10] provided an upper bound of $\lambda(G)$ is $\Delta^{2}+2 \Delta$ for a general graph with the maximum degree $\Delta$. Later, Chang and Kuo[8], improved the bound to $\Delta^{2}+\Delta$, while Kral and Skrekovski [2] reduced the bound to $\Delta^{2}+$ $\Delta-1$. Furthermore, recently Gonccalves [1] proved the bound $\Delta^{2}+\Delta-2$ which is the present best record. If $G$ is a graph of diameter 2 then $\lambda(G) \leq \Delta^{2}$. The upper bound is attainable for Moore graphs (diameter 2 graphs with order $\Delta^{2}+1$ ). (Such graphs exist only if $\Delta=2,3,7$ and possibly 57). Thus Griggs and Yeh [10] conjectured that the best bound is $\Delta^{2}$ for any graph $G$ with the maximum degree $\Delta \geq 2$. (This is not true for $\Delta=1$. For example, $\lambda\left(K_{2}\right)=1$ but $\lambda\left(K_{2}\right)=2$ ).

Graph products play an important role in connecting various useful networks and they also serve as natural tools for different concepts in many areas of research. In this paper, we have considered the graph formed by the $\gamma$ product of graphs [6] and obtained a general upper bound for $L(2,1)$-labeling number in terms of the maximum degrees of the graphs. In the case of $\gamma$-product of graphs, $L(2,1)$-labeling number of graph holds Griggs and Yeh's conjecture [10] with minor exception.

## 2. A LABELING ALGORITHM

A subset $X$ of $V(G)$ is called an $i$-stable set (or $i$ independent set) if the distance between any two vertices in $X$ is greater than $i$. An 1 -stable (independent) set is a usual independent set. A maximal 2 -stable subset $X$ of a set $Y$ is a 2 -stable subset of $Y$ such that $X$ is not a proper subset of any 2 -stable subset of $Y$.

Chang and Kuo [8] proposed the following algorithm to obtain an $L(2,1)$-labeling and the maximum value of that labeling on a given graph.
Algorithm
Input: A graph $G=(V, E)$
Output: The value $k$ is the maximum label.
Idea: In each step $i$, find a maximal 2 -stable set from the unlabeled vertices that are distance at least two away from those vertices labeled in the previous step. Then label all
the vertices in that 2 -stable set with the index $i$ in the current stage. The label $i$ starts from 0 and then increase by 1 in each step. The maximum label $k$ is the final value of $i$.

Initialization: Set $X_{-1}=\phi ; V=V(G) ; i=0$.

## Iteration:

1. Determine $Y_{i}$ and $X_{i}$.

■ $Y_{i}=\left\{u \in V: u\right.$ is unlabelled and $\left.d(u, v) \geq 2 \forall v \in X_{i-1}\right\}$.

- $X_{i}$ is a maximal 2 -stable subset of $Y_{i}$.
- If $Y_{i}=\phi$ then $\operatorname{set} X_{i}=\phi$.

2. Label the vertices of $X_{i}$ (if there is any) with $i$.
3. $V \leftarrow V-X_{i}$.
4. If $V \neq \phi$, then $i \leftarrow i+1$, go to step 1 .
5. Record the current $i$ as $k$ (which is the maximum label). Stop.

Thus $k$ is an upper bound on $\lambda(G)$. Let $u$ be a vertex with largest label $k$ obtained by above algorithm. Set

$$
I_{1}=\left\{i: 0 \leq i \leq k-1 \text { and } d(u, v)=1 \text { for some } v \in X_{i}\right\}
$$

$I_{2}=\{i: 0 \leq i \leq k-1 \quad$ and $d(u, v) \leq 2$ for some $\left.v \in X_{i}\right\}$.

$$
I_{3}=\left\{i: 0 \leq i \leq k-1 \text { and } d(u, v) \geq 3 \text { for all } v \in X_{i}\right\} .
$$

Then Chang and Kuo showed that $\lambda(G) \leq k \leq\left|I_{2}\right|+$ $\left|I_{3}\right| \leq\left|I_{2}\right|+\left|I_{1}\right|$.
In order to find $k$, it suffices to estimate $B=\left|I_{1}\right|+\left|I_{2}\right|$ in term of $\Delta(G)$. We will investigate the value $B$ with respect to a particular graph ( $\gamma$-product of two graphs). The notations which have been introduced in this section will also be used in the following sections.

## 3. THE $\boldsymbol{\gamma}$-PRODUCT OF GRAPHS

The $\gamma$-product $G \square H$ of two graphs $G$ and $H$ is the graph with vertex set $V(G) \times V(H)$, in which the vertex $(u, v)$ is adjacent to the vertex $\left(u^{\prime}, v^{\prime}\right)$ if and only if either $u$ is
adjacent to $u^{\prime}$ in $G$ or $v$ is adjacent to $v^{\prime}$ in $H$. For example, we consider the Fig.1.

By the definition of the $\gamma$-product of two graphs $G$ and $H$, if $\Delta(G)=0$ or $\Delta(H)=0$ then $G \square H$ consists of disjoint copies of $H$ or $G$. Thus $\lambda(G \boxtimes H)=\lambda(H)$ or $\lambda(G \square H)=$ $\lambda(G)$. Therefore we assume that $\Delta(G) \geq 1$ and $\Delta(H) \geq 1$.

## 4. UPPER BOUND FOR THE $L(2,1)-$ LABELING NUMBER IN G $\downarrow$ H

In this section, general upper bound for the $L(2,1)$ labeling number ( $\lambda$-number) of $\gamma$-product $G \square H$ in term of maximum degree of the graphs has been established. In this regard, we state and prove the following theorem.

Theorem 4.1. Let $\Delta, \Delta_{1}, \Delta_{2}$ be the maximum degree of $G \boxtimes H, G, H$ and $n, n_{1}, n_{2}$ be the number of vertices of $G \boxtimes H, G, H$ respectively. Then

$$
\begin{aligned}
\lambda(G \boxtimes H) \leq \Delta^{2}- & n_{1} \Delta_{2}\left(\Delta_{1}-1\right)-n_{2} \Delta_{1}\left(\Delta_{2}-1\right) \\
& -\Delta_{1} \Delta_{2}\left(\Delta_{1}+\Delta_{2}+1\right)+1 .
\end{aligned}
$$

Proof: Let $u_{\gamma}=(u, v)$ be any vertex in the graph $G \square H$. Denote $d=\operatorname{deg}_{G \boxminus H}\left(u_{\gamma}\right), d_{1}=\operatorname{deg}_{G}(u), d_{2}=\operatorname{deg}_{H}(v)$, $\Delta_{1}=\operatorname{maxdeg}(G), \Delta_{2}=\operatorname{maxdeg}(H),|V(G)|=n_{1}$ and $|V(H)|=n_{2}$. Then by the definition of $\gamma$-product we have the following results $d=n_{1} d_{2}+n_{2} d_{1}-d_{1} d_{2}$ and $\Delta=$ $n_{1} \Delta_{2}+n_{2} \Delta_{1}-\Delta_{1} \Delta_{2}$.

Let us consider the Fig.2. For any vertex $v^{\prime}$ in $H$ with distance 2 from $v$, there must be a path $v^{\prime} v^{\prime \prime} v$ of length two between $v^{\prime}$ and $v$ in $H$; but the degree of $u$ in $G$ is $d_{1}$ i.e. $u$ has $d_{1}$ adjacent vertices in $G$, by the definition of $\gamma$ product $G$ Q, there must be $d_{1}+1$ internally-disjoint paths(two paths are said to be internally-disjoint if they do not intersect each other) of length two between ( $u, v^{\prime}$ ) and $(u, v)$. Hence for any vertex in $H$ with distance 2 from $v$, there must be corresponding $d_{1}+1$ vertices with distance 2 from $u_{\gamma}=(u, v)$ which are coincided in $G \boxtimes H$; on the contrary whenever there is no such vertex in $H$ with distance 2 from $v$ in $H$, the corresponding $d_{1}+1$ vertices with distance 2 from the vertex $u_{\gamma}=(u, v)$ which are


Fig. $1 \gamma$-product of graphs
coincided in $G \boxtimes H$ will never exit. In the former case since such $d_{1}+1$ vertices with distance 2 from $u_{\gamma}=$ $(u, v)$ coincide in $G \boxtimes H$ and hence they can only be counted once and therefore we have to deduct $d_{1}+1-$ $1=d_{1}$ from the value $d(\Delta-1)$ which is best possible number of vertices at distance 2 from a vertex $u_{\gamma}=$ $(u, v)$ in $G \square H$. Let the number of vertices in $H$ with distance 2 from $v$ be $t$, then $t \in\left[0, d_{2}\left(\Delta_{2}-1\right)\right]$. If we take $t=d_{2}\left(\Delta_{2}-1\right)$ which is best possible number of vertices at distance 2 from a vertex $v$ in $H$, then to get the number of vertices at distance 2 from $u_{\gamma}=(u, v)$ in $G \boxtimes H$, we will have to subtract at least $d_{2}\left(\Delta_{2}-1\right) d_{1}$ from the value $d(\Delta-1)$. For $G$, we can proceed in the similar way to get the number of vertices at distance 2 from $u_{\gamma}=(u, v)$ in $G \boxtimes H$ and in this case subtract $d_{1}\left(\Delta_{1}-1\right) d_{2}$ from the value $d(\Delta-1)$. Hence the number of vertices at distance 2 from $u_{\gamma}=(u, v)$ in $G H$ will decrease $d_{1}\left(\Delta_{1}-1\right) d_{2}+d_{2}\left(\Delta_{2}-1\right) d_{1}$ from the value $d(\Delta-1)$ altogether. By the above analysis, the number $d(\Delta-1)-d_{1}\left(\Delta_{1}-1\right) d_{2}-d_{2}\left(\Delta_{2}-1\right) d_{1}$ is now the best possible number of vertices at distance 2 from $u_{r}=(u, v)$ in $G \boxtimes H$.

Moreover by the definition of $\gamma$-product $G \boxminus H$, we can again analyse as follows:

Let $\varepsilon$ be the number of edges of the subgraph $F$ induced by the neighbours of $u_{\gamma}$. The edges of the subgraph $F$ induced by the neighbours of $u_{\gamma}$ can be divided into the following two cases.

Case I: Consider the Fig. 3 for this case. For each neighbour vertex $\left(u, v^{\prime}\right)$ (where $v^{\prime}$ is adjacent to $v$ in $H$ ) of $u_{\gamma}=(u, v)$ and any vertex $\left(u^{\prime}, v_{t}\right)$ (where $u^{\prime}$ is adjacent to $u$ in $G$ and $v_{t}$ is any vertex of $\left.H\right),\left(u^{\prime}, v_{t}\right)$ must be the common neighbour of $\left(u, v^{\prime}\right)$ and $(u, v)$, then there must be an edge between $\left(u^{\prime}, v_{t}\right)$ and $\left(u, v^{\prime}\right)$ and an edge between $\left(u^{\prime}, v_{t}\right)$ and $(u, v)$ respectively. But there are at least $n_{2} d_{1}$ neighbour vertices like $\left(u^{\prime}, v_{t}\right)$ of $u_{\gamma}=(u, v)$ and there are totally $d_{2}$ neighbour vertices like $\left(u, v^{\prime}\right)$ of $u_{\gamma}=(u, v)$. Hence the number of edges of the subgraph $F$ induced by the neighbours of $u_{\gamma}$ is at least $n_{2} d_{1} d_{2}$ i.e. $\varepsilon \geq n_{2} d_{1} d_{2}$. By a symmetric analysis, the neighbours of $u_{\gamma}=(u, v)$ should again add at least $\left(n_{1} d_{1} d_{2}-1\right) \quad$ (excluding the coincided edge between $\left(u^{\prime}, v\right)$ and $\left.\left(u, v^{\prime}\right)\right)$.


Fig 2


Fig 3


Fig 4
Case II: Consider the Fig. 4 for this case. If $u^{\prime}$ is adjacent to $u$ in $G$, then ( $u, v$ ) must be adjacent to $\left(u^{\prime}, v_{1}\right)$ and $\left(u^{\prime}, v_{2}\right)$ where $v_{1}$ and $v_{2}$ are any two vertices of $H$, hence the vertices of the subgraph $F$ induced by the neighbours of $u_{r}=(u, v)$ should be all $\left(u^{\prime}, v\right)$ where $v \in V(H)$. Because $\Delta(H)=\Delta_{2}$ and there are totally $d_{1}$ neighbour vertices $u^{\prime}$ of $u$, then the number of edges of the subgraph $F$ induced by the neighbours of $u_{\gamma}=(u, v)$ should be greater than $d_{1} \Delta_{2}$. Hence, at least $d_{1} \Delta_{2}$ should be added to the number of edges of the subgraph $F$ induced by the neighbours of $u_{\gamma}=(u, v)$. By a symmetric analysis, the number of edges of the subgraph $F$ induced by the neighbours of $u_{\gamma}=(u, v)$ must be increased by the number $d_{2} \Delta_{1}$ at least.

By the analysis of the above two cases, we have $\varepsilon \geq$ $n_{1} d_{1} d_{2}-1+n_{2} d_{1} d_{2}+d_{1} \Delta_{2}+d_{2} \Delta_{1}$.
Whenever there is an edge in $F$, the number of vertices with distance 2 from $u_{\gamma}$ will decrease by 2 , hence the number of vertices with distance 2 from $u_{\gamma}=(u, v)$ in $G \boxtimes H$ will still need at least a decrease $n_{1} d_{1} d_{2}-1+$ $n_{2} d_{1} d_{2}+d_{1} \Delta_{2}+d_{2} \Delta_{1}$ from the value $d(\Delta-1)-$ $d_{1}\left(\Delta_{1}-1\right) d_{2}-d_{2}\left(\Delta_{2}-1\right) d_{1}$. (The number $d(\Delta-1)-$ $d_{1}\left(\Delta_{1}-1\right) d_{2}-d_{2}\left(\Delta_{2}-1\right) d_{1}$ is now the best possible for the number of vertices with distance 2 from $u_{\gamma}=$ $(u, v)$ in $G$ 凹 $)$.

Hence for any vertex $u_{\gamma}$, the number of vertices with distance 1 from $u_{\gamma}$ is no greater than $\Delta$. The number of vertices with distance 2 from $u_{\gamma}$ is no greater than
$d(\Delta-1)-d_{1}\left(\Delta_{1}-1\right) d_{2}-d_{2}\left(\Delta_{2}-1\right) d_{1}-n_{1} d_{1} d_{2}+$ $1-n_{2} d_{1} d_{2}-d_{1} \Delta_{2}-d_{2} \Delta_{1}$.

Hence $\left|I_{1}\right| \leq d .\left|I_{2}\right| \leq d+d(\Delta-1)-d_{1}\left(\Delta_{1}-1\right) d_{2}-$ $d_{2}\left(\Delta_{2}-1\right) d_{1}-n_{1} d_{1} d_{2}+1-n_{2} d_{1} d_{2}-d_{1} \Delta_{2}-d_{2} \Delta_{1}$.
Then $\quad B=\left|I_{1}\right|+\left|I_{2}\right| \leq d+d \Delta-d_{1}\left(\Delta_{1}-1\right) d_{2}-$ $d_{2}\left(\Delta_{2}-1\right) d_{1}-n_{1} d_{1} d_{2}+1-n_{2} d_{1} d_{2}-d_{1} \Delta_{2}-d_{2} \Delta_{1}=$ $\left(n_{1} d_{2}+n_{2} d_{1}-d_{1} d_{2}\right)\left(n_{1} \Delta_{2}+n_{2} \Delta_{1}-\Delta_{1} \Delta_{2}+1\right)-$ $d_{1}\left(\Delta_{1}-1\right) d_{2}-d_{2}\left(\Delta_{2}-1\right) d_{1}-n_{1} d_{1} d_{2}+1-$ $n_{2} d_{1} d_{2}-d_{1} \Delta_{2}-d_{2} \Delta_{1}$.

Define $\quad f(s, t)=\left(n_{1} t+n_{2} s-s t\right)\left(n_{1} \Delta_{2}+n_{2} \Delta_{1}-\right.$ $\left.\Delta_{1} \Delta_{2}+1\right)-s\left(\Delta_{1}-1\right) t-t\left(\Delta_{2}-1\right) s-n_{1} s t+1-$ $n_{2} s t-s \Delta_{2}-t \Delta_{1}$.

Then $f(s, t)$ has the absolute maximum at $\left(\Delta_{1}, \Delta_{2}\right)$ on $\left[0, \Delta_{1}\right] \times\left[0, \Delta_{2}\right]$.

$$
\begin{aligned}
f\left(\Delta_{1}, \Delta_{2}\right)=\left(n_{1} \Delta_{2}\right. & +n_{2} \Delta_{1} \\
& \left.-\Delta_{1} \Delta_{2}\right)\left(n_{1} \Delta_{2}+n_{2} \Delta_{1}-\Delta_{1} \Delta_{2}+1\right) \\
& -\Delta_{1}\left(\Delta_{1}-1\right) \Delta_{2}-\Delta_{2}\left(\Delta_{2}-1\right) \Delta_{1} \\
& -n_{1} \Delta_{1} \Delta_{2}+1-n_{2} \Delta_{1} \Delta_{2}-\Delta_{1} \Delta_{2} \\
& -\Delta_{2} \Delta_{1} \\
=\Delta(\Delta+ & 1)-\left(n_{1}+n_{2}+\Delta_{1}+\Delta_{2}\right) \Delta_{1} \Delta_{2}+1 \\
& =\Delta^{2}+\left(n_{1} \Delta_{2}+n_{2} \Delta_{1}-\Delta_{1} \Delta_{2}\right) \\
& -\left(n_{1}+n_{2}+\Delta_{1}+\Delta_{2}\right) \Delta_{1} \Delta_{2} \\
& +1 \\
& =\Delta^{2}-n_{1} \Delta_{2}\left(\Delta_{1}-1\right) \\
& -n_{2} \Delta_{1}\left(\Delta_{2}-1\right) \\
& -\Delta_{1} \Delta_{2}\left(\Delta_{1}+\Delta_{2}+1\right)+1 .
\end{aligned}
$$

The $\quad \lambda(G \square H) \leq k \leq B \leq \Delta^{2}-n_{1} \Delta_{2}\left(\Delta_{1}-1\right)-$ $n_{2} \Delta_{1}\left(\Delta_{2}-1\right)-\Delta_{1} \Delta_{2}\left(\Delta_{1}+\Delta_{2}+1\right)+1$.

## 5. IMPROVED BOUND FOR THE L(2,1)-LABELING NUMBER IN G $\cdot \mathbf{H}$

In this section, we shall improve the upper bound obtained in theorem 4.1 of the $L(2,1)$-labeling number on the $\gamma$ product $G \square H$ of two graphs $G$ and $H$ on the analysis of the adjacency matrices of the graph involved.

Suppose $A_{1}$ and $A_{2}$ are the adjacency matrices of $G$ and $H$ respectively. Then the adjacency matrix of $\gamma$-product $G \square$ $H$ of the graphs $G$ and $H$ with $(\bmod 2)$ can be written as $A=\left(A_{1} \otimes A_{2}\right)+\left(J_{1} \otimes A_{2}\right)+\left(A_{1} \otimes J_{2}\right)$ where $J_{1}$ is the square matrix of order $n_{1}$ with all entries 1 and $J_{2}$ is the square matrix of order $n_{2}$ with all entries 1 . These matrices involve the Kronecker product $\otimes$ of matrices, $\left(A_{1} \otimes A_{2}\right)$ is the Kronecker product $\otimes$ of matrices $A_{1}$ and $A_{2}$. Similarly $\left(J_{1} \otimes A_{2}\right)$ and $\left(A_{1} \otimes J_{2}\right)$ are Kronecker product of matrices involved in it. Note that the rules of algebra of Kronecker product $\otimes$ of matrices can be found in [5].
In order to find $k$, it suffices to estimate $B=\left|I_{1}\right|+\left|I_{2}\right|$ in term of $\Delta(G)$ (using labeling algorithm). Before eliminating the upper bound $k$, we introduce a notation first. Let $M$ be a matrix with $n$ rows. For $1 \leq i \leq n, r_{i}(M)$ denote the number of nonzero entries in the $i$ th row of $M$ excluding the diagonal entry.

Let $A$ be the adjacency matrix of $G$ with respect to a list of vertices $\left\{v_{1}, v_{2}, \ldots \ldots, v_{n}\right\}$. Then it is well known that the $(i, j)$ th entry of $A^{k}$ is the number of different $\left(v_{i}, v_{j}\right)$ walks in $G$ of length $k$, for $k \geq 0$.
Thus $r_{i}(A)=\operatorname{deg}\left(v_{i}\right), r_{i}\left(A^{2}\right)$ is the number of vertices joining by a walk of length 2 from $v_{i}$ excluding $v_{i}$ itself and $r_{i}\left(A^{2}+A\right)$ is the number of vertices of distance 1 or 2 from $v_{i}$. So that

$$
\begin{align*}
& r_{i}\left(A^{2}\right) \leq \operatorname{deg}\left(v_{i}\right)(\Delta(G)-1)  \tag{1}\\
& r_{i}\left(A^{2}+A\right) \leq \operatorname{deg}\left(v_{i}\right) \Delta(G) \tag{2}
\end{align*}
$$

For convenience, the notations which have been introduced in this section will also be used in the following section.

## 6. MAIN RESULT

Theorem 5.1: Let $\Delta_{1}, \Delta_{2}$ be the maximum degree of $G, H$ and $n_{1}, n_{2}$ be the number of vertices of $G, H$ respectively. Then
$\lambda(G \boxtimes H) \leq \Delta_{1}^{2} \Delta_{2}^{2}-\Delta_{1}^{2} \Delta_{2}-\Delta_{2}^{2} \Delta_{1}+\left(n_{2}-1\right) \Delta_{1}^{2}+$ $\left(n_{1}-1\right) \Delta_{2}^{2}+\Delta_{1} \Delta_{2}+n_{1} \Delta_{2}+n_{2} \Delta_{1}$.

Proof: From the above discussion in section 5, we get that the adjacency matrix of $G \backsim H$ is $A=\left(A_{1} \otimes A_{2}\right)+$ $\left(J_{1} \otimes A_{2}\right)+\left(A_{1} \otimes J_{2}\right)$. Then

$$
\begin{aligned}
& A^{2}+A=\left[\left(A_{1} \otimes A_{2}\right)+\left(J_{1} \otimes A_{2}\right)+\left(A_{1} \otimes J_{2}\right)\right]^{2} \\
&+\left(A_{1} \otimes A_{2}\right)+\left(J_{1} \otimes A_{2}\right)+\left(A_{1} \otimes J_{2}\right) \\
&=\left(A_{1}^{2} \otimes A_{2}^{2}\right)+\left(J_{1}^{2} \otimes A_{2}^{2}\right)+\left(A_{1}^{2} \otimes J_{2}^{2}\right) \\
&+2\left(J_{1} \otimes A_{2}\right)\left(A_{1} \otimes J_{2}\right)+2\left(A_{1} J_{1} \otimes A_{2}^{2}\right) \\
&+2\left(A_{1}^{2} \otimes A_{2} J_{2}\right)+\left(A_{1} \otimes A_{2}\right) \\
&+\left(J_{1} \otimes A_{2}\right)+\left(A_{1} \otimes J_{2}\right) \\
&=\left(A_{1}^{2} \otimes A_{2}^{2}\right)+n_{1}\left(J_{1} \otimes A_{2}^{2}\right)+ \\
& n_{2}\left(A_{1}^{2} \otimes J_{2}\right)+2\left(J_{1} A_{1} \otimes A_{2} J_{2}\right)+ \\
& 2\left(A_{1} J_{1} \otimes A_{2}^{2}\right)+2\left(A_{1}^{2} \otimes A_{2} J_{2}\right)+ \\
&\left(A_{1} \otimes A_{2}\right)+\left(J_{1} \otimes A_{2}\right)+\left(A_{1} \otimes J_{2}\right)
\end{aligned}
$$

Note that the rules of algebra of Kronecker product $\otimes$ of matrices can be found in [5].

Since all entries of the involved matrices are nonnegative, then the number of non-zero entries in the $\left(u_{i}, v_{j}\right)$ th entry of $n_{2}\left(A_{1}^{2} \otimes J_{2}\right)+2\left(A_{1}^{2} \otimes A_{2} J_{2}\right)+\left(A_{1} \otimes J_{2}\right)+n_{1}\left(J_{1} \otimes A_{2}^{2}\right)+$ $2\left(A_{1} J_{1} \otimes A_{2}^{2}\right)+\left(J_{1} \otimes A_{2}\right)+2\left(J_{1} A_{1} \otimes A_{2} J_{2}\right)$ is the same as that of $\left(A_{1}^{2} \otimes J_{2}\right)+\left(A_{1} \otimes J_{2}\right)+\left(J_{1} \otimes A_{2}^{2}\right)+\left(J_{1} \otimes A_{2}\right)=$ $\left(A_{1}^{2}+A_{1}\right) \otimes J_{2}+J_{1} \otimes\left(A_{2}^{2}+A_{2}\right)$.

Let $k$ be the maximum label obtained by the algorithm (in section 2). Let $\left(u_{i}, v_{j}\right) \in V(G) \times V(H)$ be the vertex with the label $k$. We look at the $\left(u_{i}, v_{j}\right)$ th row of the matrix $A^{2}+A$. We have

$$
\left.\begin{array}{l}
r_{\left(u_{i}, v_{j}\right)}\left(A^{2}+A\right) \leq \\
\quad r_{\left(u_{i}, v_{j}\right)}\left(A_{1}^{2} \otimes A_{2}^{2}\right) \\
\\
+r_{\left(u_{i}, v_{j}\right)}\left(\left(A_{1}^{2}+A_{1}\right) \otimes J_{2}\right) \\
\\
+r_{\left(u_{i}, v_{j}\right)}\left(J_{1} \otimes\left(A_{2}^{2}+A_{2}\right)\right) \\
\\
+r_{\left(u_{i}, v_{j}\right)}\left(A_{1} \otimes A_{2}\right)
\end{array}\right\} \begin{aligned}
=r_{i}\left(A_{1}^{2}\right) r_{j}\left(A_{2}^{2}\right)+r_{i}\left(A_{1}^{2}+A_{1}\right) r_{j}\left(J_{2}\right)+r_{i}\left(J_{1}\right) r_{j}\left(A_{2}^{2}+\right. \\
\left.A_{2}\right)+r_{i}\left(A_{1}\right) r_{j}\left(A_{2}\right)=\operatorname{deg}_{G}\left(u_{i}\right)\left(\Delta_{1}-1\right) \operatorname{deg}_{H}\left(v_{j}\right)\left(\Delta_{2}-\right. \\
1)+\operatorname{deg}_{G}\left(u_{i}\right) \Delta_{1}\left(n_{2}-1\right)+\left(n_{1}-1\right) \operatorname{deg}_{H}\left(v_{j}\right) \Delta_{2}+ \\
\operatorname{deg}_{G}\left(u_{i}\right) \operatorname{deg}_{H}\left(v_{j}\right) .
\end{aligned}
$$

Note that the last equality is obtained by applying equation (1) and (2).

Thus, the number of non-zero entries in the $\left(u_{i}, v_{j}\right)$ th entry of $\left(A^{2}+A\right)$ excluding the diagonal entry is at most $\operatorname{deg}_{G}\left(u_{i}\right)\left(\Delta_{1}-1\right) \operatorname{deg}_{H}\left(v_{j}\right)\left(\Delta_{2}-1\right)+$ $\operatorname{deg}_{G}\left(u_{i}\right) \Delta_{1}\left(n_{2}-1\right)+\left(n_{1}-1\right) \operatorname{deg}_{H}\left(v_{j}\right) \Delta_{2}+$ $\operatorname{deg}_{G}\left(u_{i}\right) \operatorname{deg}_{H}\left(v_{j}\right)$. Also we have known that $\left|I_{1}\right| \leq$ $\Delta(G \boxtimes H)=n_{1} \Delta_{2}+n_{2} \Delta_{1}-\Delta_{1} \Delta_{2}$.

Thus $\lambda(G \boxtimes H) \leq\left|I_{2}\right|+\left|I_{1}\right| \leq \Delta_{1}\left(\Delta_{1}-1\right) \Delta_{2}\left(\Delta_{2}-1\right)+$ $\Delta_{1} \Delta_{1}\left(n_{2}-1\right)+\left(n_{1}-1\right) \Delta_{2} \Delta_{2}+\Delta_{1} \Delta_{2}+n_{1} \Delta_{2}+n_{2} \Delta_{1}-$ $\Delta_{1} \Delta_{2}$.
Hence $\quad \lambda(G \boxtimes H) \leq \Delta_{1}^{2} \Delta_{2}^{2}-\Delta_{1}^{2} \Delta_{2}-\Delta_{2}^{2} \Delta_{1}+$ $\left(n_{2}-1\right) \Delta_{1}^{2}+\left(n_{1}-1\right) \Delta_{2}^{2}+\Delta_{1} \Delta_{2}+n_{1} \Delta_{2}+n_{2} \Delta_{1}$.

This completes the proof.

## 7. CONCLUSION

In theorem 4.1, we have proved that $\lambda(G \boxtimes H) \leq \Delta^{2}-$ $n_{1} \Delta_{2}\left(\Delta_{1}-1\right)-n_{2} \Delta_{1}\left(\Delta_{2}-1\right)-\Delta_{1} \Delta_{2}\left(\Delta_{1}+\Delta_{2}+1\right)+1$, where the maximum degree of $G \square H$ is $n_{1} \Delta_{2}+n_{2} \Delta_{1}-$ $\Delta_{1} \Delta_{2}$. Since $\left(\Delta^{2}-n_{1} \Delta_{2}\left(\Delta_{1}-1\right)-n_{2} \Delta_{1}\left(\Delta_{2}-1\right)-\right.$ $\left.\Delta_{1} \Delta_{2}\left(\Delta_{1}+\Delta_{2}+1\right)+1\right)-\left(\Delta_{1}^{2} \Delta_{2}^{2}-\Delta_{1}^{2} \Delta_{2}-\Delta_{2}^{2} \Delta_{1}+\right.$ $\left.\left(n_{2}-1\right) \Delta_{1}^{2}+\left(n_{1}-1\right) \Delta_{2}^{2}+\Delta_{1} \Delta_{2}+n_{1} \Delta_{2}+n_{2} \Delta_{1}\right)=$ $\left(n_{1} \Delta_{2}+n_{2} \Delta_{1}\right)^{2}-\left(n_{2}-1\right) \Delta_{1}^{2}-\left(n_{1}-1\right) \Delta_{2}^{2}-$
$\Delta_{1} \Delta_{2}\left(2 n_{1} \Delta_{2}+2 n_{2} \Delta_{1}+n_{1}+n_{2}+2\right)+1$. We have thus reduced the bound by $\left(n_{1} \Delta_{2}+n_{2} \Delta_{1}\right)^{2}-\left(n_{2}-1\right) \Delta_{1}^{2}-$ $\left(n_{1}-1\right) \Delta_{2}^{2}-\Delta_{1} \Delta_{2}\left(2 n_{1} \Delta_{2}+2 n_{2} \Delta_{1}+n_{1}+n_{2}+2\right)+1$.

## 8. ACKNOWLEDGMENTS

This research work is supported by University Grant commission (UGC) New Delhi, India under the UGC National Fellowship for OBC student scheme to the first author.

## 9. REFERENCES

[1] D.Gonccalves, On the $L(p, 1)$-labeling of graphs, in Proc. 2005 Eur. Conf. Combinatorics, Graph Theory Appl. S. Felsner, Ed., (2005), 81-86 .
[2] D. Kral and R. Skrekovski, A theorem about channel assignment problem, SIAM J. Discrete Math., 16 (2003) 426-437.
[3] D. D. F. Liu and R. K. Yeh, On Distance Two Labeling of Graphs, Ars Combinatoria, 47 (1997) 1322.
[4] D. Sakai, Labeling Chordal Graphs with a condition at distance two, SIAM J. Discrete Math., 7 (1994) 133-140.
[5] D. M. Cvetkovic, M. Doob, and H. Sachs, Spectra of graphs-Theory and application, $2^{\text {nd }}$ ed. New York: Academic, 1982.
[6] E. M. El-Kholy, E .S. Lashin and S. N. Daoud, New operations on graphs and graph foldings, International Mathematical Forum, 7(2012) 2253-2268.
[7] F. S .Roberts, "T-colorings of graphs; Recent results and open problems," Discr. Math. vol 93, pp. 229245, 1991.
[8] G. J. Chang and D. Kuo, The L(2, 1) -labeling on graphs, SIAM J. Discrete Math., 9 (1996) 309-316.
[9] G. J. Chang and et al., On L(d, 1) -labeling of graphs, Discrete Math, 220 (2000) 57-66.
[10] J. R. Griggs and R. K. Yeh, Labeling graphs with a condition at distance two, SIAM J. Discrete Math., 5 (1992) 586-595.
[11] P. K. Jha, Optimal $L(2,1)$-labeling of strong product of cycles, IEEE Trans. Circuits systems-I, Fundam. Theory Appl., 48(4) (2001) 498-500.
[12] P. K. Jha, Optimal L(2, 1) -labeling on Cartesian products of cycles with an application to independent
domination, IEEE Trans. Circuits systems-I, Fundam. Theory Appl., 47(12) (2000) 1531-1534.
[13] P. Pradhan, K. Kumar, The L(2,1)-labeling on $\alpha$ product of Graphs, Annals of Pure and Applied Mathematics, Vol. 10, No. 1, 2015, 29-39.
[14] S. Klavzar and S. Spacepan, The $\Delta^{2}$-conjecture for $\mathrm{L}(2,1)$-labelings is true for direct and strong products of graphs, IEEE Trans. Circuits systems-II, Exp. Briefs, 53(3) (2006) 274-277.
[15] S. K. Vaidya and D. D. Bantva, Distance two labeling of some total graphs, Gen. Math. Notes, 3(1) (2011) 100-107.
[16] S. K. Vaidya and D. D. Bantva, Some new perspectives on distance two labeling, International Journal of Mathematics and Soft Computing, 3(3) (2013), 7-13.
[17] S. Paul, M. Pal and A. Pal, L(2,1) -labelling of circular-arc graph, Annals of Pure and Applied Mathematics, 5(2) (2014) 208-219.
[18] S. Paul, M. Pal and A. Pal, L(2,1)-labeling of permutation and bipartite permutation graphs, Mathematics in Computer Science, DOI 10.1007/s11786-014-0180-2.
[19] S. Paul, M. Pal and A. Pal, L(2,1) -labeling of interval graphs, J. Appl. Math. Comput. DOI 10.1007/s12190-014-0846-6.
[20] W. K. Hale, Frequency assignment: Theory and application, Proc. IEEE, 68(6) (1980) 1497-1514.
[21] W. C. Shiu, Z. Shao, K. K. Poon and D. Zhang, A new approach to the $L(2,1)$-Labeling of some products of graphs, IEEE Trans. Circuits and systemsII, Express Briefs, Vol. 55. No. 8 August (2008) 1549-7747.
[22] Z. Shao and et al., Improved bounds on the $L(2,1)$ number of direct and strong products of graphs, IEEE Trans. Circuits systems-II, Exp. Briefs, 55(7) (2008) 685-689.
[23] Z. Shao and R .K. Yeh, The L(2, 1) -labeling and operations of graphs, IEEE Trans. Circuits and systems-I, 52(4) (2005) 668-671.
[24] Z. Shao and D. Zhang, The $L(2,1)$-labeling on Cartesian sum of graphs, Applied Mathematics Letters, 21 (2008) 843-848.
[25] Z. Shao and et al., The $L(2,1)$-labeling of $K_{1, n}$-free graphs and its applications, Applied Mathematics Letters, 21 (2008) 1188-1193.
[26] Z. Shao and et al., The $L(2,1)$-labeling on graphs and the frequency assignment problem, Applied Mathematics Letters, 21 (2008) 37-41.

