

# Can one find External Source Input Expressions for which there exist Map Reduce Configurations?

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## ABSTRACT

An intention of MapReduce Sets for External Source Input expressions analysis has to suggest criteria how External Source Input expressions in External Source Input data can be defined in a meaningful way and how they should be compared. Similitude based MapReduce Sets for External Source Input Expression Analysis and MapReduce Sets for Assignment is expected to adhere to fundamental principles of the scientific External Source Input process that are expressiveness of External Source Input models and reproducibility of their External Source Input inference. External Source Input expressions are assumed to be elements of a External Source Input expression space or Conjecture class and External Source Input data provide “information” which of these External Source Input expressions should be used to interpret the External Source Input data. An inference External Source Input algorithm constructs the mapping between External Source Input data and External Source Input expressions, in particular by a External Source Input cost minimization process. Fluctuations in the External Source Input data often limit the External Source Input precision, which we can achieve to uniquely identify a single External Source Input expression as interpretation of the External Source Input data. We advocate an information theoretic perspective on External Source Input expression analysis to resolve this dilemma where the tradeoff between External Source Input informativeness of statistical inference External Source Input and their External Source Input stability is mirrored in the information-theoretic External Source Input optimum of high External Source Input information rate and zero communication expression error. The inference External Source Input algorithm is considered as an outlier object External Source Input path, which naturally limits the resolution of the External Source Input expression space given the uncertainty of the External Source Input data.

## Keywords

MapReduce, External Source Input expressions, kernel function.

## 1. INTRODUCTION

The field of general expression intention is to formalize the notion of expressions in precise mathematical terms [1]. Expressions are perceived as regular structures behind the External Source Input data sources, i.e., “the underlying deep regular structures are descriptions of the source, which are hidden via the sensing path” [3]. Expression Intention combines algebra, geometry and statistics to explain the nature of External Source Input data sources and, thereby, depicts a generative External Source Input modeling perspective on expression analysis. This philosophy argues for a distinct generative viewpoint to infer the probability

distribution of the External Source Input data. In many real-world situations the External Source Input data are represented or generated in a very high dimensional space and the information-processing task focuses on a low dimensional interpretation space. The analysis of visual External Source Input data like images or videos provides a very convincing example of this situation: intensity expressions that are sensed by a camera are mathematically represented as points in a space with  $\#\{\text{intensities}\}^{\#\{\text{pixels}\}}$  dimensions. When segmenting an image in semantically distinct regions then the interpretation space contains  $\#\{\text{segments}\}^{\#\{\text{sites}\}}$  elements where the number of sites is often much smaller than the number of pixels. The reader should note that the space of segmentations is still exponentially large in the number of sites. Consequently, we adopt a discriminative view of MapReduce Sets for External Source Input Expression assignment: the External Source Input expressions which are inferred from the External Source Input data are elements of an interpretation space called Conjecture class and these External Source Input expressions are more or less closely related with the External Source Input data generating mechanism of the source. The Conjecture class often also reflects information about the aim of expression analysis, i.e., what the External Source Input expressions are used for in subsequent information processing.

The inference of External Source Input expressions in External Source Input data is formulated as an algorithmic search for a stable subset of the underlying Conjecture class. Stability is required to guarantee that the External Source Input expression analysis process would yield an equivalent outcome for the same structure of the External Source Input data source but a different realization of the outlier object process. A second, antagonistic requirement of the External Source Input expression analysis process is its specificity or informativeness: a small subset of the Conjecture class and in the limit a single Conjecture should be selected which poses a tradeoff to the stability requirement. Both External Source Input principles mirror the reproducibility and specificity requirement of the scientific reasoning [9, 15].

External Source Input expression analysis algorithms often follow an optimization principle. Desired External Source Input expressions are assigned a high score or low costs and undesirable External Source Input expressions are discarded by assigning a low score or high costs. In the following, we adopt the terminology of cost minimization rather than score maximization. A kernel function defines a partial order of Conjecture minimal kernels distinguish where the most preferred Conjecture. The outlier object in the External Source Input data, however, may introduce fluctuations in the kernels and the Conjecture with minimal kernels for one realization of the External Source Input data may no longer

minimize kernels for a second realization of the External Source Input data. Therefore, we advocate stabilizing the set of kernel-minimal Conjecture by expanding it to a set of Conjecture with near-optimal kernels, also called approximation set. The size of such an approximation set is determined by information theoretic considerations. Conjecture in the approximation set is considered to be statistically indistinguishable.

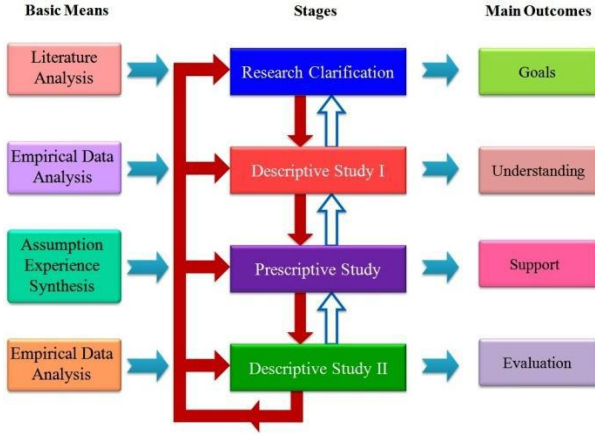


Fig 1: Research Plan: Basic Means, Stages, Main Outcomes

## 2. RESEARCH CLARIFICATION

External Source Input expression analysis quantifies structures in External Source Input data, which usually relate to a set of objects. To mathematically characterize this problem domain we have to define what we mean by measurements and Conjecture. Given is a set of objects  $\mathbf{O}^{(n)} = \{O_1, \dots, O_n\} \subset \mathcal{O}$ ,  $n \in \mathbb{N}$ . Individual objects can be characterized by measurements either relative to an external reference frame, e.g., a coordinate system in a feature space, or by comparison to other objects. A measurement  $X$  is defined as a mapping of an object configuration in a measurement space, i.e.,

$$X: \mathcal{O}^1 \times \dots \times \mathcal{O}^r \rightarrow \mathbb{K}, \quad (O_1, \dots, O_r) \mapsto X_{O_1, \dots, O_r}. \quad (1)$$

The object configurations are often specified as collections of objects taken from the same object set  $\mathcal{O}^1 = \dots = \mathcal{O}^r$ . For binary External Source Input data, the first and the second object set can differ  $\mathcal{O}^1 \neq \mathcal{O}^2$ . The most often used measurement types are feature vectors  $X: \mathcal{O} \rightarrow \mathbb{R}^d$  denoted as  $X_o \in \mathbb{R}^d$ . Relational External Source Input data arise often in informatics applications and in network analysis problems. They are defined as  $X: \mathcal{O} \times \mathcal{O} \rightarrow \mathbb{R}$ , where  $X_{o_1, o_2}$  denotes a proximity/Similitude value between object  $O_1$  and  $O_2$ . More complicated External Source Input data structures than vectors or relations, e.g., three-way External Source Input data or (hyper) graphs, are employed in various applications.

In the following, we use the generic notation  $\mathbf{X}^{(n)} \in \mathcal{X}^{(n)}$  for a set of measurements to characterize these  $n$  objects  $\mathbf{O}^{(n)}$ .  $\mathcal{X}^{(n)}$  denotes the corresponding measurement space of  $n$  objects. To simplify notation we omit the index  $^{(n)}$  whenever the dependence on problem size is clear.

A Conjecture  $c(\cdot)$  of a MapReduce Sets for External Source Input Expression assignment problem is a function that assigns a set of objects or a set of object configurations to a External Source Input expression out of a External Source Input expression space  $\mathcal{C}$ ; i.e.,

$$c: \mathcal{O}^1 \times \dots \times \mathcal{O}^r \rightarrow \mathcal{C}, \quad (O_1, \dots, O_r) \mapsto c(O_1, \dots, O_r). \quad (2)$$

The intention of Conjecture does not depend on the measurements  $X_{O_1, \dots, O_r}$  but potential External Source Input expressions that are denoted by Conjecture are defined prior to any measurements. The reader should note that the notion of a “feasible solution” in applied mathematics and optimization often depends on constraints that are determined by measurements contrary to the intention in (2). Such situations can be modeled by unconstrained solution spaces with infinite kernels for those solutions that violate the constraints.

The Conjecture class for a MapReduce Sets for External Source Input Expression assignment problem is defined as the set of functions assigning an object or an object configuration (In the following, we restrict Conjecture to map an object to a External Source Input expression. The more general situation of object configurations can be analyzed in an analogous way but involves a more complex notation) to an element of the External Source Input expression space, i.e.,

$$\mathcal{C}(\mathbf{O}) = \{c(\mathbf{O}) : \mathbf{O} \in \mathcal{O}\}. \quad (3)$$

A well-known example of a Conjecture class is the space of partitions or classification functions  $c: \mathcal{O} \rightarrow \{1, \dots, k\}$  which we use in classification or clustering. When clustering  $n$  objects into  $k$  clusters, then we restrict the space of all possible partition functions to  $\mathcal{C}^{(n)} = \{1, \dots, k\}^n$  for the object set  $\mathbf{O}^{(n)}$ . The corresponding Conjecture class is denoted by  $\mathcal{C}^{(n)} = \mathcal{C}(\mathbf{O}^{(n)})$ . For parameter estimation problems the External Source Input expressions are possible values of the matrices and the External Source Input expression space is a subset of the  $d$ -dimensional Euclidean rotations.

The Conjecture class is a set of functions that map objects or object configurations to expressions. External Source Input expression analysis requires assessing the quality of Conjecture  $c \in \mathcal{C}$ . We adopt a kernel function viewpoint in this paper, which attributes a non-negative kernel value

$$R: \mathcal{C}^{(n)} \times \mathcal{X}^{(n)} \rightarrow \mathbb{R}_+, \quad (c, \mathbf{X}^{(n)}) \mapsto R(c, \mathbf{X}^{(n)}) \quad (4)$$

to each Conjecture given the measurements ( $\mathbb{R}_+ := [0, \infty)$ ). The non-negativity assumption does not restrict the choice of kernel functions since we can always replace  $\tilde{R}(c, \mathbf{X}^{(n)}) := R(c, \mathbf{X}^{(n)}) - \inf_{c \in \mathcal{C}} R(c, \mathbf{X}^{(n)})$  for effectively computable minimal kernels.

The classical intention of statistical Inference [2] advocates to use the posteriori minimize as the solution of the inference problem. The best posteriori External Source Input expression denoted by  $c^\perp(\mathbf{X}^{(n)})$  minimizes the posteriori loss function of the External Source Input expression analysis problem given the measurements  $\mathbf{X}^{(n)}$ , i.e.,

$$c^\perp(\mathbf{X}^{(n)}) \in \arg \min_{c \in \mathcal{C}^{(n)}} R(c, \mathbf{X}^{(n)}). \quad (5)$$

Although Conjecture map objects into a External Source Input expression space, the posteriori loss function minimizer  $c^\perp(\mathbf{X}^{(n)})$  depends on measurements.

The intention requires for inferability of classifications that the Conjecture class is not “too complex” (i.e., finite computing) and, as a consequence, the solution  $c^\perp(\mathbf{X}^{(n)})$  converges to the optimal solution which minimizes the expected loss function. A corresponding criterion has been derived for regression [4].

This classical inference intention is not applicable when the size of the Conjecture class grows with the number of objects

like in clustering or other optimization problems of combinatorial nature. Without strong regularization we cannot hope to identify a single solution, which globally minimizes the expected loss function in the asymptotic limit  $n \rightarrow \infty$ . Conjecture class of combinatorial problems often have an infinite computing and, therefore, are not inferable in the classical computing sense. Therefore, we replace the concept of a unique function as the solution of an inference problem with a weighted set of kernel functions. The challenge of inference then amounts to determine a weight measure, which is concentrated on few solutions to achieve precision. The weights  $w$  are defined as functions which map triplets of a Conjecture, measurements and a resolution parameter to the unit interval, i.e.,

$$w: \mathcal{C}^{(n)} \times \mathcal{D}^{(n)} \times \mathbb{R}_+ \rightarrow [0,1], \quad (c, \mathbf{X}^{(n)}, \beta) \mapsto w_\beta(c, \mathbf{X}^{(n)}). \quad (6)$$

The set of weights is denoted as

$$\mathcal{W}_\beta(\mathbf{X}^{(n)}) = \{w_\beta(c, \mathbf{X}^{(n)}) : c \in \mathcal{C}^{(n)}\}.$$

How should we choose the weights  $w_\beta(c, \mathbf{X}^{(n)})$  that large weights are only assigned to kernel functions with low kernels? The partial ordering constraint

$$R(c, \mathbf{X}^{(n)}) \leq R(\tilde{c}, \mathbf{X}^{(n)}) \Leftrightarrow w_\beta(c, \mathbf{X}^{(n)}) \geq w_\beta(\tilde{c}, \mathbf{X}^{(n)}), \quad (7)$$

ensures that kernel functions with minimal kernels  $R(c^\perp, \mathbf{X}^{(n)})$  assume the maximal weight value. Weights are normalized to one i.e.,  $0 \leq w_\beta(c, \mathbf{X}^{(n)}) \leq 1$ . The non-negativity constraint of weights allows us to write the weights as  $w_\beta(c, \mathbf{X}^{(n)}) = \exp(-\beta f(R(c, \mathbf{X}^{(n)})))$  with the monotonic kernel function  $f(x)$ . Since  $f(x)$  amounts to a monotone rescaling of the kernels  $R(c, \mathbf{X}^{(n)})$  we resort to the common choice of Statistical weights with the inverse computational value  $\beta$ , i.e.,

$$w_\beta(c, \mathbf{X}^{(n)}) = \exp(-\beta R(c, \mathbf{X}^{(n)})). \quad (8)$$

It is worth mentioning that standard approximation sets as introduced in the intention of approximation algorithms would correspond to binary weights

$$w_\beta^{bin}(c, \mathbf{X}^{(n)}) = \begin{cases} 1 & \text{if } R(c, \mathbf{X}^{(n)}) \leq R(c^\perp, \mathbf{X}^{(n)}) + 1/\beta \\ 0 & \text{otherwise} \end{cases} \quad (9)$$

The weight  $w_\beta(c, \mathbf{X}^{(n)})$  of a given Conjecture  $c$  is a random variable of the measurements  $\mathbf{X}^{(n)}$ . We consider the quantity

$$Z_\beta(\mathbf{X}^{(n)}) := \sum_{c \in \mathcal{C}^{(n)}} w_\beta(c, \mathbf{X}^{(n)}), \quad (10)$$

which measures the total weight of Conjecture with low kernels. The weight sum is also known as the partition function in statistics when we use Statistical weights. In case of binary weights  $Z_\beta(\mathbf{X}^{(n)})$  denotes the number of solutions that are  $1/\beta$  close to the optimum.

To determine the optimal regularization of a MapReduce Sets for External Source Input Expression assignment method we have to define and estimate the generalization performance of Conjecture. We adopt the two instance scenario with training and test External Source Input data described by respective object sets  $\mathbf{O}'$ ,  $\mathbf{O}''$  and corresponding measurements  $\mathbf{X}'$ ,  $\mathbf{X}'' \sim \mathbb{P}(\mathbf{X})$ . Both sets of measurements are drawn Independent and identically distributed from the same probability distribution  $\mathbb{P}(\mathbf{X})$ . The training and test External Source Input data  $\mathbf{X}'$ ,  $\mathbf{X}''$  define two optimization problems  $R(\cdot, \mathbf{X}')$ ,  $R(\cdot, \mathbf{X}'')$ . The two instance scenario or two sample set scenario is widely used in statistics and statistical Inference intention [5], i.e., to bound the deviation of posteriori loss function from expected loss function, but also for two-terminal systems in information intention [6, 13]. Statistical expression analysis

requires that inferred External Source Input expressions have to generalize from training External Source Input data to test External Source Input data since outlier object in the External Source Input data might render the solution  $c^\perp(\mathbf{X}') \neq c^\perp(\mathbf{X}'')$  unstable. How can we evaluate the generalization properties of solutions to a MapReduce Sets for External Source Input Expression assignment problem? Before we can compute the kernels  $R(\cdot, \mathbf{X}'')$  on test External Source Input data of approximate solutions  $c(\mathbf{O}') \in \mathcal{C}(\mathbf{O}')$  on training External Source Input data we have to identify a External Source Input expression  $c(\mathbf{O}'') \in \mathcal{C}(\mathbf{O}'')$  which corresponds to  $c(\mathbf{O}')$ . A priori, it is not clear how to compare External Source Input expressions  $c(\mathbf{O}')$  for objects  $\mathbf{O}'$  with External Source Input expressions  $c(\mathbf{O}'')$  for objects  $\mathbf{O}''$ . Therefore, we define a bijective mapping

$$\psi: \mathcal{O}' \rightarrow \mathcal{O}'', \quad \mathbf{O}' \mapsto \psi \circ \mathbf{O}'. \quad (11)$$

The mapping  $\psi$  allows us to identify a External Source Input expression Conjecture for training set of objects  $c' \in \mathcal{C}(\mathbf{O}')$  with a External Source Input expression Conjecture for a test set of objects  $c'' \in \mathcal{C}(\psi \circ \mathbf{O}')$ . The reader should note that such a mapping  $\psi$  might change the object indices. In cases when the objects  $\mathbf{O}'$ ,  $\mathbf{O}''$  are elements of an underlying metric space, then a natural choice for  $\psi$  is the nearest neighbor mapping. The mapping  $\psi$  enables us to evaluate External Source Input expression kernels on test External Source Input data  $\mathbf{X}''$  for External Source Input expressions  $c(\mathbf{O}')$  selected on the basis of training External Source Input data  $\mathbf{X}'$ . Consequently, we can determine how many training External Source Input expressions with large weights share also large weights on test External Source Input data, i.e.,

$$\Delta Z_\beta(\mathbf{X}', \mathbf{X}'') := \sum_{c \in \mathcal{C}(\mathbf{O}')} w_\beta(c, \psi \circ \mathbf{X}') w_\beta(c, \mathbf{X}''). \quad (12)$$

A large subset of Conjecture with jointly large weights indicates that low kernel Conjecture on training External Source Input data  $\mathbf{X}'$  also perform with low kernels on test External Source Input data. The tradeoff between stability and informativeness for Statistical weights on (8) is controlled by maximizing  $\beta$  for given loss function  $R(\cdot, \mathbf{X})$  under the constraint of large weight overlap  $\Delta Z_\beta(\mathbf{X}', \mathbf{X}'') / Z_\beta(\mathbf{X}'') \approx$

1. A quantitative statement how close this ratio should approach unity requires a statistical decision intention as provided by information transmission.

A natural question in statistical inference arises from asymptotic considerations in the large  $n$ -limit. What is the asymptotic behavior of the log weight sum  $\log Z_\beta(\mathbf{X}^{(n)})$  dependent on the problem/instance size  $n$ ? As remarked above the measurements  $\mathbf{X}^{(n)}$  of a particular MapReduce Sets for External Source Input Expression assignment instance depend on the value  $n$ . In analogy to information intention [7, 14] we assume that the log weight sums converge according to an asymptotic equipartition property, i.e.,

$$\mathcal{F}' := \lim_{n \rightarrow \infty} - \frac{\log Z_\beta(\mathbf{X}'^{(n)})}{\log |\mathcal{C}(\mathbf{O}'^{(n)})|}, \quad (13)$$

$$\mathcal{F}'' := \lim_{n \rightarrow \infty} - \frac{\log Z_\beta(\mathbf{X}''^{(n)})}{\log |\mathcal{C}(\mathbf{O}''^{(n)})|}, \quad (14)$$

$$\Delta_{\mathcal{F}} := \lim_{n \rightarrow \infty} - \frac{\log \Delta Z_{\beta}(\mathbf{X}'^{(n)}, \mathbf{X}''^{(n)})}{\log |\mathcal{C}(\mathbf{O}''^{(n)})|}. \quad (15)$$

These assumptions (13, 14, 15) requires that the log weight sums normalized by the size of the Conjecture class converge towards deterministic limits. The quantities  $\mathcal{F}'$ ,  $\mathcal{F}''$  are known as the key values (up to a factor  $\beta^{-1}$ ) for the instances  $R(\cdot, \mathbf{X}')$ ,  $R(\cdot, \mathbf{X}'')$  in statistically. The factor  $\log |\mathcal{C}^{(n)}|$  denotes the problem size of the optimization problem, i.e., it is  $O(n)$  for clustering problems with maximally  $k^n$  different partitions and  $O(n \log n)$  for sorting problems with  $\log |\mathcal{C}^{(n)}| = \log(n!)$ .

**Intention 1.** The set  $A_{\varepsilon}^{(n)}$  of jointly typical instances w.r.t.  $p(\mathbf{X}'^{(n)}, \mathbf{X}''^{(n)})$  is the set of instance pairs  $(\mathbf{X}'^{(n)}, \mathbf{X}''^{(n)}) \in \mathcal{G}^{(n)} \times \mathcal{G}^{(n)}$  with posteriori log partition functions close to the respective key values

$$A_{\varepsilon}^{(n)} = \{(\mathbf{X}'^{(n)}, \mathbf{X}''^{(n)}) \in \mathcal{G}^{(n)} \times \mathcal{G}^{(n)} : \left| - \frac{\log Z_{\beta}(\mathbf{X}'^{(n)})}{\log |\mathcal{C}(\mathbf{O}'^{(n)})|} - \mathcal{F}' \right| < \varepsilon, \quad (16)$$

$$\left| - \frac{\log Z_{\beta}(\mathbf{X}''^{(n)})}{\log |\mathcal{C}(\mathbf{O}''^{(n)})|} - \mathcal{F}'' \right| < \varepsilon, \quad (17)$$

$$\left| - \frac{\log \Delta Z_{\beta}(\mathbf{X}'^{(n)}, \mathbf{X}''^{(n)})}{\log |\mathcal{C}(\mathbf{O}''^{(n)})|} - \Delta_{\mathcal{F}} \right| < \varepsilon \}. \quad (18)$$

The reader should note that the weak law of large numbers guarantees convergence of posteriori entropies towards their expectation values in information intention. Due to the dependence of the weights  $w_{\beta}(c, \mathbf{X}^{(n)})$  on the kernel function  $R(\cdot, \mathbf{X}^{(n)})$  convergence has to be required for a kernel function. We also conjecture that kernel functions that violate this convergence behavior cannot be used to define predictive External Source Input models.

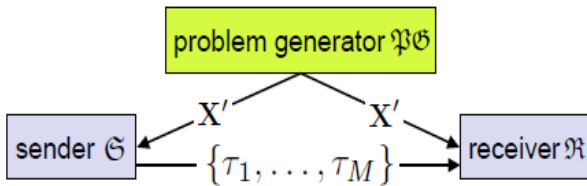


Fig 2: Generation of a set of  $M$  instruction problems by e.g. permuting the object indices.

### 3. DESCRIPTIVE STUDY I

In the following, we describe an information theoretic framework to determine which Conjecture are statistically indistinguishable due to outlier object in the measurements and how much we have to coarsen the Conjecture class. Random instruction concept suggests a External Source Input model intention to determine the maximal number of distinguishable  $n$ -key value strings in the Group of Hamiltonian space when the key value strings are exposed to outlier object in a communication path. We develop a generalization of this idea for solution spaces of optimization problems. The weight distribution  $w_{\beta}(c, \mathbf{X}^{(n)})$ ,  $c \in \mathcal{C}$  over the Conjecture class  $\mathcal{C}$  corresponds to the subsets of key value strings assigned to a specific instruction vector in information intention. Outlier object perturbs the measurements and

therefore, the weight distribution fluctuates. An algorithm to approximately minimize a kernel function and the measurements as input to this algorithm defines an outlier object path in an asymptotical communication scenario with a mapper  $M$ , a reducer  $R$  and a key generator  $KG$ [12, 18]. The key generator connects the mapper with the reducer by posing an optimization problem given a kernel function or an algorithm. Communication takes place by approximately optimizing a given kernel function, i.e., by calculating weight sets  $Z_{\beta}(\mathbf{X}')$ ,  $Z_{\beta}(\mathbf{X}'')$ . This instruction concept will be referred to as approximation set instruction (ASI) since the weights are concentrated on approximate minimizers of the optimization problem. The outlier object path is characterized by an External Source Input expression kernel function  $R(c, \mathbf{X})$  that determines the path capacity of the ASI scenario. Selection and validation of MapReduce Sets for External Source Input Expression assignment models are then achieved by maximizing the path capacity over a set of kernel functions  $R_{\theta}(\cdot, \mathbf{X})$ ,  $\theta \in \Theta$  where  $\theta$  indexes the various kernel functions or MapReduce Sets for External Source Input Expression assignment objectives. In a more general setting an arbitrary algorithm which does not necessarily minimize a kernel function can be considered to define a weight distribution and thereby, to play the role of an outlier object path [8,16] due to fluctuations in the input or in the execution path.

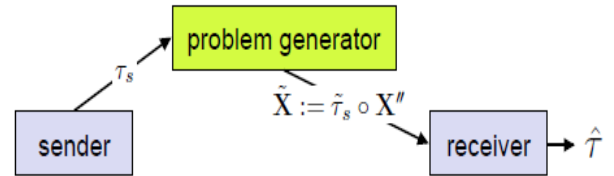


Fig 3: Communication process: (1) the mapper selects transformation  $\tau_s$ , (2) the key generator draws  $\mathbf{X}'' \sim \mathbb{P}(\mathbf{X})$  and applies  $\tilde{\tau}_s = \psi \circ \tau_s \circ \psi^{-1}$  to it, and the reducer estimates  $\hat{t}$  based on  $\tilde{\mathbf{X}} = \tilde{\tau}_s \circ \mathbf{X}''$ .

Before we describe the communication protocol we have to define the instruction for communication. We introduce random instruction intention to demonstrate the limits of asymptotically error free communication over an outlier object path. Random instruction refers to the fact that messages in random instruction External Source Input model are selected as a set of key value strings  $\{\xi^{(j)} = (\xi_1^{(j)}, \dots, \xi_n^{(j)}), 1 \leq j \leq M\}$  with length  $n = |\log |\mathcal{C}(\mathbf{O}^{(n)})|$  that are drawn Independent and identically distributed according to a probability distribution  $p(\xi)$ . For sufficiently large  $n$ , the instructions all have mutual variations which are highly concentrated around the expected variation  $2np(1-p)$  with the probability  $p = \mathbb{P}(\xi^{(1)} = 1)$ . In the asymptotic limit  $n \rightarrow \infty$  for  $p = 1/2$ , the random instructions uniformly partition the Group of Hamiltonian space of  $n$ -key value sequences into subsets of key value strings which can be reverse instruction without errors. In an analogous way, we cover the Conjecture class by weight distributions [10, 17]. To generate a uniform cover of the Conjecture class, we introduce a transformation

$$\tau : \mathcal{C} \rightarrow \mathcal{C}, \quad \mathbf{O} \mapsto \tau \circ \mathbf{O}. \quad (19)$$

The set of all possible transformations is denoted as  $\mathbb{T}$ . Transformations that are restricted to object sets  $\mathbf{O}^{(n)}$  of  $n$  objects are denoted by  $\tau^{(n)} \in \mathbb{T}^{(n)}$ . A random cover of the Conjecture class is then generated by selecting a set of transformations  $\mathcal{T} = \{\tau_j^{(n)} \in \mathbb{T}^{(n)} : 1 \leq j \leq M, \tau_j^{(n)} \sim P(\tau^{(n)})\}$  with a rate  $\rho := \log M / \log |\mathbb{T}^{(n)}|$ . A natural choice of the probability distribution for transformations is the uniform distribution  $P(\tau^{(n)}) = 1/|\mathbb{T}^{(n)}|$ . The intuition behind the

transformations is the following: When a transformation is applied to an object set  $\mathbf{O}$  then the respective Conjecture  $c(\mathbf{O})$  and the measurements  $X_{\mathbf{O}}$  are transformed accordingly. Furthermore, the weights  $w_{\beta}(c, \mathbf{X})$  are transformed by applying  $\tau$  to  $c$  and  $\mathbf{X}$ , i.e.,  $\tau \circ w_{\beta}(c, \mathbf{X}) := w_{\beta}(\tau \circ c, \tau \circ \mathbf{X})$ .

Analogous to random instruction strategy, we generate the transformations  $\tau^{(n)} \sim P(\tau^{(n)})$  in a random way. The probability distribution  $P(\tau^{(n)})$  is defined over the set of possible transformations  $\mathbb{T}^{(n)}$ . An asymptotic equipartition property depends on the entropy density of the transformation set

$$\mathcal{H}(\tau) := \lim_{n \rightarrow \infty} - \frac{\log P(\tau^{(n)})}{\log |\mathbb{T}^{(n)}|}. \quad (20)$$

For instruction, we choose  $\varepsilon$ -typical transformations  $\tau^{(n)} \in \mathbb{T}_{\varepsilon}^{(n)}$  with the typical set  $\mathbb{T}_{\varepsilon}^{(n)}$  being defined in the following way:

**Intention 2.** The set  $\mathbb{T}_{\varepsilon}^{(n)}$  of typical transformations w.r.t.  $P(\tau^{(n)})$  is the set of transformations  $\tau^{(n)} \in \mathbb{T}^{(n)}$  with the property

$$\mathbb{T}_{\varepsilon}^{(n)} = \{ \tau^{(n)} \in \mathbb{T}^{(n)} : \left| - \frac{\log P(\tau^{(n)})}{\log |\mathbb{T}^{(n)}|} - \mathcal{H}(\tau) \right| < \varepsilon \}. \quad (21)$$

Special cases of such transformations  $\tilde{\tau}^{(n)}$  are random permutations when optimizing combinatorial optimization kernel functions like clustering External Source Input models or graph cut problems. In parametric statistics, the transformations are parameter grids of e.g. rotations when estimating the orthogonal transformations of PCA or SVD.

Mapper  $\mathfrak{S}$  and reducer  $\mathfrak{R}$  agree on a kernel function for MapReduce Sets for External Source Input Expression assignment  $R(c, \mathbf{X}')$  and on a mapping function  $\psi$ . The following procedure is then employed to generate the instruction for the communication process:

1. Mapper  $\mathfrak{S}$  and reducer  $\mathfrak{R}$  obtain External Source Input data  $\mathbf{X}'$  from the key generator  $\mathfrak{B}\mathfrak{G}$ .
2.  $\mathfrak{S}$  and  $\mathfrak{R}$  calculate the weight set  $\mathcal{O}_{\beta}(\mathbf{X}')$ .
3.  $\mathfrak{S}$  generates a set of (random) transformations  $\mathcal{T} := \{\tau_1, \dots, \tau_{2^{np}}\}$ . The transformations define a set of optimization problems  $R(c, \tau_j \circ \mathbf{X}')$ ,  $1 \leq j \leq 2^{np}$  to determine weight sets  $\mathcal{O}_{\beta}(\tau_j \circ \mathbf{X}')$ ,  $1 \leq j \leq 2^{np}$ .
4.  $\mathfrak{S}$  sends the set of transformations  $\mathcal{T}$  to  $\mathfrak{R}$  who determines the set of weight sets  $\{\mathcal{O}_{\beta}(\tau_j \circ \mathbf{X}')\}_{j=1}^{2^{np}}$ .

The reason behind this procedure is the following: Given the measurements  $\mathbf{X}'$  the mapper has randomly covered the Conjecture class  $\mathcal{C}(\mathbf{O}')$  by respective weight sets  $\{\mathcal{O}_{\beta}(\tau_j \circ \mathbf{X}') : 1 \leq j \leq 2^{np}\}$ . Communication can take place if the weight sets are stable under the stochastic fluctuations of the measurements. The criterion for reliable communication is defined by the ability of the reducer to identify the transformation which has been selected by the mapper. After this setup procedure, both mapper and reducer have a list of weight sets available.

How is the communication between mapper and reducer organized? During communication, the following steps take place as depicted in fig. 3:

1. The mapper  $\mathfrak{S}$  selects a transformation  $\tau_s$  as message and sends it to the key generator  $\mathfrak{B}\mathfrak{G}$ .
2.  $\mathfrak{B}\mathfrak{G}$  generates a new External Source Input data set  $\mathbf{X}''$  and

establishes correspondence  $\psi$  between  $\mathbf{X}'$  and  $\mathbf{X}''$ .  $\mathfrak{B}\mathfrak{G}$  then applies the selected transformation  $\tau_s$ , yielding  $\tilde{\mathbf{X}} = \psi \circ \tau_s \circ \psi^{-1} \circ \mathbf{X}''$ .

3.  $\mathfrak{B}\mathfrak{G}$  sends  $\tilde{\mathbf{X}}$  to the reducer  $\mathfrak{R}$  without revealing  $\tau_s$ .

4.  $\mathfrak{R}$  calculates the weight set  $\mathcal{O}_{\beta}(\tilde{\mathbf{X}})$ .

5.  $\mathfrak{R}$  estimates the selected transformation  $\tau_s$ , by using reverse instruction rule

$$\hat{\tau} = \arg \max_{\tau \in \mathcal{T}} \sum_{c \in \mathcal{C}(\mathbf{O}')} w_{\beta}(c, \psi \circ \tau \circ \mathbf{X}') w_{\beta}(c, \tilde{\mathbf{X}}). \quad (22)$$

In the case of discrete Conjecture class, then the communication path is bounded from above by the cardinality of  $\mathcal{C}(\mathbf{X})$  if two conditions hold: (i) the path is outlier object free  $\mathbf{X}' \equiv \mathbf{X}''$ ; (ii) the transformation set is sufficiently rich that every Conjecture can be selected as a global minimizer of the kernel function.

#### 4. PRESCRIPTIVE STUDY I

To determine the optimal approximation precision for an optimization problem  $R(\cdot, \mathbf{X})$  we have to derive necessary and sufficient conditions which have to hold in order to reliably identify the transformations  $\tau_s \in \mathcal{T}$ . The parameter  $\beta$ , which controls the concentration of weights and thereby the resolution of the Conjecture class, has to be adapted to the size of the transformation set  $|\mathcal{T}|$ . Therefore, we analyze the error probability of the reverse instruction rule (22) which is associated with a particular kernel function  $R(\cdot, \mathbf{X})$  a rate  $\rho$ . The maximal value of  $\beta$  under the condition of zero error communication is defined as *approximation capacity* since it determines the approximation precision of the instruction scheme.

A communication error occurs if the mapper selects  $\tau_s$  and the reducer reverse instructions  $\hat{\tau} = \tau_j$ ,  $j \neq s$ . To estimate the probability of this event, we introduce the weight overlaps

$$\Delta Z_{\beta}^j := \sum_{c \in \mathcal{C}(\mathbf{O}')} w_{\beta}(c, \psi \circ \tau_j \circ \mathbf{X}') w_{\beta}(c, \tilde{\mathbf{X}}), \tau_j \in \mathcal{T}. \quad (23)$$

The quantity  $\Delta Z_{\beta}^j$  measures the number of Conjecture, which have jointly low kernels  $R(c, \psi \circ \tau_j \circ \mathbf{X}')$  and  $R(c, \tilde{\mathbf{X}})$ .

The probability of a communication error is given by a substantial overlap  $\Delta Z_{\beta}^j$  induced by  $\tau_j \in \mathcal{T} \setminus \{\tau_s\}$ ,  $1 \leq j \leq M$ ,  $j \neq s$ , i.e.,

$$\begin{aligned} \mathbb{P}(\hat{\tau} \neq \tau_s | \tau_s) &= \mathbb{P}(\max_{1 \leq j \leq M, j \neq s} \Delta Z_{\beta}^j \geq \Delta Z_{\beta}^s | \tau_s) \\ &\stackrel{(a)}{\leq} \sum_{1 \leq j \leq M, j \neq s} \mathbb{P}(\Delta Z_{\beta}^j \geq \Delta Z_{\beta}^s | \tau_s) \\ &\stackrel{(b)}{\leq} \sum_{1 \leq j \leq M, j \neq s} \mathbb{E}_{\mathbf{X}', \mathbf{X}''} \mathbb{E}_{\mathcal{T} \setminus \{\tau_s\}} \left[ \frac{\Delta Z_{\beta}^j}{\Delta Z_{\beta}^s} \middle| \tau_s \right] \\ &\stackrel{(c)}{=} (M-1) \mathbb{E}_{\mathbf{X}', \mathbf{X}''} \left[ \frac{\mathbb{E}_{\tau_j: j \neq s} [\Delta Z_{\beta}^j | \mathbf{X}', \mathbf{X}'']}{\Delta Z_{\beta}^s} \middle| \tau_s \right] \end{aligned} \quad (24)$$

The expectation  $\mathbb{E}_{\mathcal{T} \setminus \{\tau_s\}}$  is calculated w.r.t., the set of random transformations  $\tau_j$ ,  $1 \leq j \leq M$ ,  $j \neq s$  where we have conditioned on the mapper selected transformation  $\tau_s$ . The joint probability distribution of all transformations  $\mathbb{P}(\mathcal{T}) = \prod_{j=1}^M \mathbb{P}(\tau_j)$  decomposes into product form since all transformations are randomly drawn from the set of all possible transformations  $\{\tau_j\}$ . It corresponds to the random instruction External Source Input in information intention.

The inequality (a) results from the union bound and (b) is due

to Markov's inequality. The identity (c) exploits the fact that the transformations  $\tau \in \mathcal{T}$  are Independent and identically distributed drawn according to the measure  $\mathbb{P}(\mathcal{T})$ .

The expected overlap  $\mathbb{E}_{\tau_j} \Delta Z_\beta^j$ ,  $j \neq s$  with any other message  $\tau_j$ ,  $j \neq s$  for given training External Source Input data  $\mathbf{X}'$  and test External Source Input data  $\mathbf{X}''$  conditioned on  $\tau_s$  is defined by

$$\begin{aligned} \mathbb{E}_{\tau_j, j \neq s} [\Delta Z_\beta^j | \mathbf{X}', \mathbf{X}''] &= \sum_{\tau_j \in \mathcal{T}} P(\tau_j) \\ &\sum_{c \in \mathcal{C}(\mathbf{O}'')} w_\beta(c, \psi \circ \tau_j \circ \mathbf{X}') w_\beta(c, \tilde{\mathbf{X}}) \\ &\stackrel{(d)}{\leq} \sum_{c \in \mathcal{C}(\mathbf{O}'')} w_\beta(c, \tilde{\mathbf{X}}) \exp(-\log |\mathbb{T}| (\mathcal{H}(\tau) - \varepsilon)) \\ &\quad \underbrace{\left( \sum_{\tau_j \in \mathcal{T}} w_\beta(\tau_j^{-1} \circ \psi^{-1} \circ c, \mathbf{X}') \right)} \\ &\leq Z_\beta(\mathbf{X}') \\ &\stackrel{(e)}{\leq} \exp(-\log |\mathbb{T}| (\mathcal{H}(\tau) - \varepsilon)) Z_\beta(\mathbf{X}'') Z_\beta(\mathbf{X}') \quad (25) \end{aligned}$$

The inequality (d) results from the typically of  $P(\tau_j)$ . The last inequality (e) holds since the set  $\{\tau_j^{-1} \circ \psi^{-1} \circ c : c \in \mathcal{C}(\mathbf{O}'')\}$ ,  $\tau_j \in \mathcal{T} \subset \mathcal{C}(\mathbf{O}')$ , and extending the sum  $\sum_{\tau_j \in \mathcal{T}}$  to  $\sum_{c \in \mathcal{C}(\mathbf{O}'')}$  only adds positive terms. Effectively, the sum over a random transformation  $\tau_j$  decouples the two sums and yields a of weight sums. Inserting result (25) into equation (24) yields

$$\begin{aligned} \mathbb{P}(\hat{\tau} \neq \tau_s | \tau_s) &\leq (M-1) \exp(-\log |\mathbb{T}| (\mathcal{H}(\tau) - \varepsilon)) \\ &\mathbb{E}_{\mathbf{X}', \mathbf{X}''} Z_\beta(\mathbf{X}'') Z_\beta(\mathbf{X}') / \Delta Z_\beta^s \quad (26) \\ &\stackrel{(f)}{\leq} \exp(\log M - \log |\mathbb{T}| (\mathcal{H}(\tau) - \varepsilon) - (\mathcal{F}' - \varepsilon) \log |\mathcal{C}(\mathbf{O}')| \\ &- (\mathcal{F}'' - \varepsilon) \log |\mathcal{C}(\mathbf{O}'')| + (\Delta \mathcal{F} + \varepsilon) \log |\mathcal{C}(\mathbf{O}'')|) \\ &= \exp(\log |\mathbb{T}| \rho - \log |\mathbb{T}| (\mathcal{H}(\tau) - \varepsilon) - \log |\mathcal{C}'| \mathcal{F}' - \log |\mathcal{C}''| \\ &(\mathcal{F}'' - \Delta \mathcal{F}) + \varepsilon (\log |\mathbb{T}| + \log |\mathcal{C}'| + 2 \log |\mathcal{C}''|)), \quad (27) \end{aligned}$$

where we have introduced the rate intention  $\rho = \log M / \log |\mathbb{T}|$ . The term proportional to  $\varepsilon$  can be neglected since it becomes arbitrarily small in the limit  $\lim_{n \rightarrow \infty}$  due to the assumed asymptotic equipartition property (13, 14, 15). Often, the assumption  $|\mathcal{C}(\mathbf{O}')| = |\mathcal{C}(\mathbf{O}'')| = |\mathbb{T}|$  is justified and the bound (26) simplifies to

$$\mathbb{P}(\hat{\tau} \neq \tau_s | \tau_s) \leq \exp(-\log |\mathcal{C}'| (\mathcal{F} - \rho - 4\varepsilon)) \quad (28)$$

$$\text{with } \mathcal{F}_\beta := \mathcal{H}(\tau) + \mathcal{F}' + \mathcal{F}'' - \Delta \mathcal{F} \quad (29)$$

The quantity  $\mathcal{F}_\beta$  plays the role of the mutual information in communication. Error free communication requires  $\rho < \mathcal{F}_\beta$ , i.e., the rate  $\rho$  should not exceed  $\mathcal{H}(\tau) + \mathcal{F}' + \mathcal{F}'' - \Delta \mathcal{F}$ .

How can this upper bound (28) with quantity  $\mathcal{F}_\beta$  be interpreted? A close look at equation (26) reveals that the bound depends on the term

$$\begin{aligned} \mathbb{E}_{\mathbf{X}', \mathbf{X}''} \frac{Z_\beta(\mathbf{X}'') Z_\beta(\mathbf{X}')}{|\mathbb{T}| \Delta Z_\beta^s} &= \mathbb{E}_{\mathbf{X}', \mathbf{X}''} \exp(-\log |\mathbb{T}| / Z'_\beta \\ &- \log |\mathcal{C}''| / Z''_\beta + \log |\mathcal{C}'| / Z''_\beta) \quad (30) \end{aligned}$$

with the abbreviation  $Z'_\beta = Z_\beta(\mathbf{X}')$ ,  $Z''_\beta = Z_\beta(\mathbf{X}'')$ . The term  $\log(|\mathbb{T}| / Z'_\beta)$  counts the number of ways you can form statistically distinguishable subset of the complete transformation class  $\mathbb{T}$ , the second term  $\log(|\mathcal{C}'| / Z''_\beta)$  measure the same property on the reducer side and the last term  $\log(|\mathcal{C}''| / Z''_\beta)$  accounts for double counting of the overlap. The three terms together define mutual information

between the selected message  $\tau_s$  and the reconstructed message  $\hat{\tau}$ .

## 5. PRESCRIPTIVE STUDY II:

The analysis of the error probability suggests the following inference principle for controlling the appropriate strengths which implements a form of External Source Input model selection: the approximation precision is controlled by  $\beta$  which has to be maximized to derive more precise solutions or External Source Input expressions. For small  $\beta$  the rate  $\rho$  will be low since we resolve the space of solutions only in a coarse grained fashion. For too large  $\beta$  the error probability does not vanish which indicates confusions between  $\tau_j$ ,  $j \neq s$  and  $\tau_s$ . The optimal  $\beta$ -value is given by the largest  $\beta$  or, equivalently the highest approximation precision

$$\beta^* = \arg \max_{\beta \in [0, \infty)} \mathcal{F}_\beta(\tau_s, \hat{\tau}). \quad (31)$$

Another choice to be made in modeling is to select a suitable kernel function  $R(\cdot, \mathbf{X})$  for the MapReduce Sets for External Source Input Expression assignment problems at hand. Let us assume that a number of kernel functions  $\{R_\theta(\cdot, \mathbf{X}), \theta \in \Theta\}$  are considered as candidates. The approximation capacity  $\mathcal{F}_\beta(\tau_s, \hat{\tau})$  depends on the kernel function through the weights. Therefore, we can rank the different External Source Input models according to their  $\mathcal{F}_\beta(\tau_s, \hat{\tau})$  values. Robust and informative kernel functions yield a higher approximation capacity than simpler or more brittle External Source Input models. A quantifiable choice is to select the kernel function

$$R^*(c, \mathbf{X}) = \arg \max_{\theta \in \Theta} \mathcal{F}_\beta(\tau_s, \hat{\tau} | R_\theta). \quad (32)$$

Where both the random variables  $\tau_s$  and  $\hat{\tau}$  depend on  $R_\theta(c, \mathbf{X})$ ,  $\theta \in \Theta$ . The selection rule (32) prefers the External Source Input model which is "expressive" enough to exhibit high information content (e.g., many clusters in clustering) and, at the same time robustly resists to outlier object in the External Source Input data set. The key values, which are measured in the communication setting, are context sensitive since they refer to a Conjecture class  $\mathcal{C}(\mathbf{X})$ , i.e., how finely or coarsely functions can be resolved in  $\mathcal{C}$ .

## 6. DESCRIPTIVE STUDY II

To demonstrate the approach to regularized optimization we will apply it to an almost trivial optimization problem, i.e., minimizing the Group of Hamiltonian variation to a reference key value string  $\xi' = (\xi'_1, \xi'_2, \dots, \xi'_n) \in \{-1, 1\}^n$  of  $n$  key values. This optimization problem describes the reverse instruction step in classical communication intention. The kernel function for communication measures the difference between a key value string  $s \in \{-1, 1\}^n$  and a reference instruction  $\xi'$ , i.e.,

$$R(s, \xi') = \sum_{i=1}^n \mathbb{I}\{s_i \neq \xi'_i\} = \frac{1}{2} (n - \sum_{i=1}^n s_i \xi'_i). \quad (33)$$

The variable  $s$  has to be optimized and the posteriori minimum is  $s = \xi'$ . However,  $\xi'$  is exposed to path outlier object and, in the spirit of approximation set instruction, we should only approximate it. The weights of approximate solutions are defined by

$$\mathcal{W}_\beta(\xi') = \{ w_\beta(s, \xi') = \exp(-\frac{\beta}{2} (n - \sum_{1 \leq i \leq n} s_i \xi'_i)) \}. \quad (34)$$

The mapper uses this measurement  $\xi'$  and permutes the key values according to one of the randomly selected transformations  $\mathcal{F} := \{\tau_1, \dots, \tau_{2M}\}$ . Permutations which leave  $\xi'$  invariant are excluded. This set of randomly selected transformations generates a instruction with instruction

vectors  $\{\tau_1 \circ \xi', \dots, \tau_{2M} \circ \xi'\}$ .

During communication, a second key value string  $\xi''$  is generated by the key generator. The reducer then receives the message  $\tilde{\xi} = \tau_s \circ \xi''$  when the mapper decides to communicate with transformation  $\tau_s$ . This process defines the approximation problem  $R(s, \tilde{\xi}) = \frac{1}{2} (n - s \cdot \tilde{\xi})$  on the reducer side. Based on the External Source Input data  $\tilde{\xi}$  reducer has to estimate the transformation  $\tau_s$  which has been communicated by the mapper.

Let us assume that the probability  $\delta := \mathbb{P}(\xi'_i \neq \xi''_i)$  characterizes the communication path. Therefore, a fraction  $\delta$   $n$  key values are different between the first key value sequence  $\xi'$  and the second key value sequence  $\xi''_i$ , i.e.,  $\delta = \frac{1}{n} |\{i : \xi'_i \neq \xi''_i\}|$ .

The weight sums  $Z_\beta(\xi)$ ,  $\xi \in \{\xi', \xi''\}$  are given by

$$\begin{aligned} Z_\beta(\xi) &= \sum_{s \in \mathcal{R}(\xi)} \exp(-\frac{\beta}{2} (n - \sum_{i \leq n} s_i \xi_i)) \\ &= \exp(-\beta \frac{n}{2}) \prod_{i \leq n} \sum_{s_i \in \{-1, 1\}} \exp(\frac{\beta}{2} s_i \xi_i) \\ &= \exp(-\frac{n\beta}{2}) 2^n (\cosh \frac{\beta}{2})^n. \end{aligned} \quad (35)$$

The number of jointly approximating key value strings is determined by

$$\begin{aligned} \Delta Z_\beta &= \sum_{s \in \mathcal{R}(\xi'')} \exp(-\beta(n - \frac{1}{2} \sum_{i \leq n} s_i (\xi'_i + \xi''_i))) \\ &= \exp(-\beta n) \prod_{i \leq n} (\exp(\frac{\beta}{2} (\xi'_i + \xi''_i)) + \exp(-\frac{\beta}{2} (\xi'_i + \xi''_i))) \\ &= \exp(-\beta n) 2^n (\cosh \beta)^{n(1-\delta)}. \end{aligned} \quad (36)$$

The mutual information (30) for the special case of minimizing Group of Hamiltonian variations is determined by

$$\begin{aligned} \mathcal{J}_\beta &= \mathcal{H}(\tau) + \mathcal{F}' + \mathcal{F}'' - \Delta \mathcal{F} \\ &= \ln 2 - \lim_{n \rightarrow \infty} \frac{1}{n} (\ln Z_\beta(\xi') + \ln Z_\beta(\xi'') - \ln \Delta Z_\beta) \\ &= (1-\delta) \ln \cosh \beta - 2 \ln \cosh \frac{\beta}{2} \\ &= \ln 2 + (1-\delta) \ln \cosh \beta - \ln(\cosh \beta + 1). \end{aligned} \quad (37)$$

where we have estimated the size of the set of possible random transformations as  $|\mathbb{T}| = 2^n$ . In the case of a biased sequence with  $\pi := \mathbb{P}(\xi'_i = 1) \neq \frac{1}{2}$  the cardinality of the transformation set is  $|\mathbb{T}| = 2^{\mathcal{H}(\pi)}$  with the binary entropy  $\mathcal{H}(\pi) = -\pi \log_2 \pi - (1-\pi) \log_2(1-\pi)$ .

The optimal value for  $\beta$  determined by the maximum of  $\mathcal{J}_\beta$ , i.e.,

$$\frac{d\mathcal{J}_\beta}{d\beta} = (1-\delta) \frac{\sinh \beta}{\cosh \beta} - \frac{\sinh \beta}{\cosh \beta + 1} = 0. \quad (38)$$

$$\Rightarrow \cosh \beta = \frac{1-\delta}{\delta}, \quad \cosh \beta + 1 = \delta^{-1}. \quad (39)$$

Inserting these values into equation (37) yields

$$\begin{aligned} \mathcal{J}_\beta &= \ln 2 + (1-\delta) \ln \frac{1-\delta}{\delta} - \ln \frac{1-\delta}{\delta} - \ln \frac{1}{\delta} \\ &= \ln 2 + (1-\delta) \ln(1-\delta) - \delta \ln \delta \\ &= \ln 2 - \mathcal{H}(\delta). \end{aligned} \quad (40)$$

Equation (40) shows that optimally approximating the Group of Hamiltonian variation of key value strings by

approximation set instruction yields the path capacity of the binary symmetric path with key value error probability  $\delta$ .

## 7. CONCLUSION

MapReduce Sets for External Source Input expression analysis explores the questions how similar different External Source Input expressions are and how we should compare them. The underlying topology and metric of a Conjecture class are often chosen ad hoc in applications and usually do not derive from properties of the External Source Input data source. Approximation set instruction as a External Source Input model validation principle establishes a notion of *External Source Input expression equivalence* by considering them as statistically indistinguishable when the External Source Input expression differences cannot be exploited for instruction. External Source Input expressions with the same or similar weights are considered to be equally acceptable solutions and these weights directly depend on the objective or kernel function. To justify a natural topology and metric, we have to validate the underlying kernel function for the External Source Input expression analysis problem. The reader should realize that the assumption of a kernel function assumes a lot of information about the Conjecture class; it essentially establishes a partial order of Conjecture.

External Source Input model selection and validation requires estimating the generalization ability of External Source Input models from training to test External Source Input data. “Good” External Source Input models show a high expressiveness and they are robust w.r.t., outlier object in the External Source Input data. This tradeoff between *informativeness* and *robustness* ranks different External Source Input models when they are tested on new External Source Input data and it quantitatively describes the underfitting/overfitting dilemma. In this paper we have explored the idea to use approximation sets of clustering solutions as a communication instruction. The *approximation capacity* of a kernel function provides a selection criterion, which renders various External Source Input models comparable in terms of their respective key value rates. The number of reliably extractable key values of a External Source Input expression analysis kernel function  $R(\cdot, \mathbf{X})$  defines a “task sensitive information measure” since it only accounts for the fluctuations in the External Source Input data  $\mathbf{X}$  which actually have an influence on identifying an individual External Source Input expression or a set of External Source Input expressions.

The maximum entropy inference principle suggests that we should average over the statistically indistinguishable solutions in the optimal approximation set. Such a External Source Input model averaging strategy replaces the original kernel function with the value and, thereby, it defines a continuation method with maximal robustness. Algorithmically, maximum entropy inference can be implemented by annealing methods. The entropy naturally answers the question in many External Source Input data analysis applications, which regularization term should be used without introducing an unwanted bias. The second question, how the regularization parameter should be selected, is also answered by: Choose the parameter value that maximizes the approximation capacity! The link to robust optimization is analyzed from a theoretical computer science viewpoint.

For External Source Input model selection can be applied to all combinatorial or continuous optimization problems, which depend on outlier object External Source Input data. The

outlier object level is characterized by two sample sets  $\mathbf{X}'$ ,  $\mathbf{X}''$ . It has been posteriorly explored by External Source Input model validation problems for External Source Input model based clustering of high dimensional Gaussian distributed External Source Input data and of Boolean External Source Input data. The well-known spin glass phase of maximum likelihood estimations for Gaussian sources is identified as structure with zero information content for instruction. IT can also be used to select External Source Input models for spectral clustering. Furthermore, denoising of Boolean matrices guided by the generalization capacity of SVD suggests a cutoff rank for the SVD spectrum [11].

The reader should realize that we only require an objective or kernel  $R(\cdot, \mathbf{X})$  to define a weight distribution. Any other mechanism to arrive at such a concept of approximate solutions will serve the same purpose. In principle, this concept of measuring the generalization performance of algorithms can be applied to algorithm evaluation and also to robust algorithm External Source Input. It endows the space of algorithm with a topology since two algorithms are neighbors if their approximation sets for the same input distributions share a high overlap. Such methods to measure the robustness of algorithms to errors in the computation or in the input will be in high demand to program novel hardware that trades consumption against precision of computation. So far we are completely lacking External Source Input principles for algorithm engineering, which consider this tradeoff between usage and correctness. We are also convinced that the information theoretic analysis of algorithms will shed new light on the relation between computational complexity and statistical complexity – the two faces of complexity science whose relation is far from being understood.

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