

Cost Analysis of a Two-Unit Cold Standby System Considering Hardware, Software Failures and Inspection with Maximum Repair Time

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ABSTRACT

The present paper deals with a reliability model of a two unit cold standby system that exposed to two types of failures. In each unit there are hardware and software components work together and fail independently. there is a single server who visits the system immediately when required. The server inspects the unit in case of software failure and decides either to repair or replace it. The h/w components undergo for repair at their failure and are replaced by new one in case of arriving to maximum repair time. All time distributions are considered to be exponential. Various measures of reliability of the system are obtained using regenerative point technique. Finally simulation study is done to illustrate the results.

General Terms

Stochastic Process

Keywords

Cold Standby System, Maximum Repair Time, Inspection, Replacement, Profit Gain

1. INTRODUCTION

Recently, with technological progress most of the Industrial Systems consist of software subsystems and hardware subsystems. A lot of research [6, 7, 8, 9] that have been implemented in this region have been taken into account hardware subsystem alone. But there are a lot of complex systems, which have the two subsystems together, such as control systems. [2, 11] developed a combined reliability model for the whole system in which software and hardware components work together. [3] formulated reliability models of a computer system with independent failure of software and hardware components.

Furthermore, the continued operation and ageing of these systems gradually reduce their performance and reliability. [1, 10] proved that the reliability of a system can be increased by making replacement of the components by new one when the repair time is too long. [5, 4] discussed stochastically a two unit cold standby system with the concept of priority subject to maximum operation and maximum repair time.

The aim of this paper is to study a two unit cold standby system considering hardware, software failures and inspection with maximum repair times. Explicit expressions for the following reliability measures of the system are obtained:

1. The mean time to system failure (MTSF).
2. Pointwise and steady state availability.
3. Busy periods with repair due to hardware failures or software failures or replacement or inspection.
4. The profit gain of the system, are obtained.

Finally numerical example is presented to illustrate the theoretical results.

2. ASSUMPTIONS

1. The system consists of two identical units.
2. Initially one unit is operating and the other is in standby case (cold standby).
3. The online unit suffers two types of failure, namely, software and hardware failure.
4. There is one server man and always available.
5. If the operating unit fails with software failure the server man inspects the unit and decides either to repair with some probability or replace it .
6. When the online unit exceeds maximum repair time of hardware failure, the server man will replace it.
7. If the operating unit suffers software failure while replacing the other unit so, there is priority to repair it after replacing the other unit.
8. The distributions of all times are considered to be stochastically independent random variables, each having an exponential distribution.
9. The connected switch is perfect.
10. After repair, the unit is as good as new.

3. NOTATIONS

E_0 :	State of the system at $t = 0$.
E :	Set of regenerative states.
\bar{E} :	Set of nonregenerative states.
$F_i(t)$ ($i = 1, 2$):	Cdf of hardware and software failures respectively.
$G_i(t)$ ($i = 1, 2$):	Cdf of repair times.
$q_{ij}(t), Q_{ij}(t)$:	Pdf and Cdf of time for the system transits from regenerative state S_i to S_j .
$q_{ij}^{(k)}(t), Q_{ij}^{(k)}(t)$:	Pdf and Cdf of time for the system transits from regenerative state S_i to S_j via the non-regenerative state E .
$\Pi_i(t)$:	Cdf of time to system failure starting from state $S_i \in E$.
m_{ij} :	Contribution to mean sojourn time in state S_i , when system transits direct to S_j .
μ_i :	$\int P[\text{System sojourns in state } S_i \text{ for at least time } t] dt$.
$M_i(t)$:	$P[\text{system is up initially in state } S_i \in E \text{ is up at time } t \text{ without passing through any other regenerative state or returning to itself through one or more states } \in E]$.
$AV_i(t)$:	$P[\text{The system is up at time } t \setminus E_0 = S_i \in E]$.
$k(t), K(t)$:	Pdf and Cdf of replacment time.
$m(t), M(t)$:	Pdf and Cdf of maximum repair time.
$h(t), H(t)$:	Pdf and Cdf of inspection time.
$\theta, 1 - \theta$:	probability that the unit goes for repair due to software failure/ probability that the unit will be replaced due to software failure.
$\dot{B}_i(t)$:	$P[\text{server man is busy with repair of hardware, software failures, replacement and inspection respectively at time } t \text{ starting from state } S_i \in E, (i = 1, 2, 3, 4)]$.
$C(t)$:	The expected profit incurred in $(0, t]$.
U :	Dummy variable in Laplace transform (LT).
$*$:	Symbol for LT.
\odot :	Symbol for convolution of $f(t)$ and $g(t)$ is $f(t) \odot g(t) = \int_0^t f(x) dG(t-x)$.

3.1 Symbols for the states of the system

O/S :	Unit in normal mode / in standby mode.
r_1/r_2 :	Unit is in failure mode due to hardware failure and under repair /Unit is in failure mode due to software failure and under repair.
wr_1/wr_2 :	waiting for repair of hardware failure / waiting for repair software failure .
$repl /insp$:	Unit under replacement/ Unit under inspection.

Considering these symbols, the system can be in any one of the following states.

$S_0 = (O, S)$	$S_1 = (O, r_1)$	$S_2 = (O, insp)$
$S_3 = (O, r_2)$	$S_4 = (wr_2, r_2)$	$S_5 = (wr_1, r_2)$
$S_6 = (wr_2, repl)$	$S_7 = (O, repl)$	$S_8 = (wr_2, insp)$
$S_9 = (wr_1, r_1)$	$S_{10} = (wr_2, r_1)$	$S_{11} = (wr_1, insp)$
$S_{12} = (wr_1, repl)$		

Up states : S_0, S_1, S_2, S_3 and S_7 . Down States: $S_4, S_5, S_6, S_8, S_9, S_{10}, S_{11}$ and S_{12} . One can see that all states are regenerative states.

3.2 Transition Probabilities and Sojourn Times

It can be observed that the epoch of entry into any of the states $S_i \in E$ is regenerative point. Let $T_1 (\equiv 0), T_2, T_0, \dots$ denote the epochs at which the system enters any state $S_i \in E$ let X_n denote the state Visited at epoch T_n , i.e. just after transition at T_n . $\{X_n, T_n\}$ is a Markov renewal process with state space E and $Q_{ij}(t) = P[X_{n+1} = j, T_{n+1} - T_n \leq t \setminus X_n = i]$, is the semi Markov kernel over E

The transition probability matrix of embedded Markov chain is $P = (p_{ij}) = (Q_{ij}(\infty) = Q(\infty))$, With non zero elements.

By probabilistic arguments, the non zero elements P_{ij} are,

Transition form S_0 :

$$P_{0,1} = \frac{\lambda}{\lambda + \alpha}$$

where $P_{0,1}$ means that the software failure does not occur until time t and hardware failure occurs at time t .

All other transition probabilities can be explained in the same manner.

$$P_{0,2} = \frac{\alpha}{\lambda + \alpha}$$

$$\text{Transition form } S_1 : P_{1,0} = \frac{a}{\lambda + \alpha + a + m} \quad P_{1,9} = \frac{\lambda}{\lambda + \alpha + a + m}$$

$$P_{1,10} = \frac{\alpha}{\lambda + \alpha + a + m} \quad P_{1,7} = \frac{m}{\lambda + \alpha + a + m}$$

$$\text{Transition form } S_2 : P_{2,11} = \frac{\lambda}{\lambda + \alpha + h} \quad P_{2,8} = \frac{\alpha}{\lambda + \alpha + h} \quad P_{2,7} = (1 - \theta) \frac{h}{\lambda + \alpha + h} \quad P_{2,3} = (\theta) \frac{h}{\lambda + \alpha + h}$$

$$\text{Transition form } S_3 : P_{3,5} = \frac{\lambda}{\lambda + \alpha + b} \quad P_{3,4} = \frac{\alpha}{\lambda + \alpha + b} \quad P_{3,0} = \frac{b}{\lambda + \alpha + b}$$

$$\text{Transition form } S_4 : P_{4,3} = 1$$

$$\text{Transition form } S_5 : P_{5,1} = 1$$

$$\text{Transition form } S_6 : P_{6,3} = 1$$

$$\text{Transition form } S_7 : P_{7,0} = \frac{k}{\lambda + \alpha + k} \quad P_{7,12} = \frac{\lambda}{\lambda + \alpha + k} \quad P_{7,6} = \frac{\alpha}{\lambda + \alpha + k}$$

$$\text{Transition form } S_8 : P_{8,6} = (1 - \theta) P_{8,4} = (\theta)$$

$$\text{Transition form } S_9 : P_{9,1} = 1$$

$$\text{Transition form } S_{10} : P_{10,3} = 1$$

$$\text{Transition form } S_{11} : P_{11,5} = (\theta) P_{11,12} = (1 - \theta)$$

$$\text{Transition form } S_{12} : P_{12,1} = 1$$

We can prove that, $P_{0,1} + P_{0,2} = 1, P_{1,0} + P_{1,9} + P_{1,10} + P_{1,7} = 1, P_{2,11} + P_{2,8} + P_{2,7} + P_{2,3} = 1, P_{3,0} + P_{3,4} + P_{3,5} = 1, P_{4,3} = P_{5,1} = P_{6,3} = P_{9,1} = P_{10,3} = P_{12,1} = 1, P_{7,0} + P_{7,12} + P_{7,6} = 1, P_{8,6} + P_{8,4} = 1, P_{11,5} + P_{11,12} = 1$.

The mean sojourn time μ_i in state $S_i \in E$ are, $\mu_0 = \frac{1}{\lambda + \alpha}, \mu_1 = \frac{1}{\lambda + \alpha + a + m}, \mu_2 = \frac{1}{\lambda + \alpha + h}, \mu_3 = \frac{1}{\lambda + \alpha + b}, \mu_4 = \frac{1}{b}, \mu_5 = \frac{1}{b}, \mu_6 = \frac{1}{k}, \mu_7 = \frac{1}{\lambda + \alpha + k}, \mu_8 = \frac{1}{h}, \mu_9 = \frac{1}{a}, \mu_{10} = \frac{1}{a}, \mu_{11} = \frac{1}{h}, \mu_{12} = \frac{1}{k}$.

We can prove that, $\mu_0 = m_{01} + m_{02}, \mu_1 = m_{10} + m_{19} + m_{1,10} + m_{1,7}, \mu_2 = m_{2,11} + m_{2,8} + m_{2,7} + m_{2,3}, \mu_3 = m_{3,5} + m_{3,4} + m_{3,0}, \mu_4 = m_{4,3}, \mu_5 = m_{5,1}, \mu_6 = m_{6,3}, \mu_7 = m_{7,0} + m_{7,12} + m_{7,6}, \mu_8 = m_{8,6} + m_{8,4}, \mu_9 = m_{9,1}, \mu_{10} = m_{10,3}, \mu_{11} = m_{11,5} + m_{11,12}, \mu_{12} = m_{12,1}$.

3.3 MEAN TIME TO SYSTEM FAILURE

According to the arguments of theory of regenerative processes, we obtain the following equations,

$$\bar{\Pi}_0(t) = \bar{F}_1(t) \bar{F}_2(t) + q_{01}(t) \odot \bar{\Pi}_1(t) + q_{02}(t) \odot \bar{\Pi}_2(t) \quad (1)$$

The above equation can be explained as follows: The left hand side means that the system works until time t . The first term in the right hand side means that the hardware failure does not occur until time t and software failure not occur until time t . The second term means that the system moves from state S_0 to state S_1 at a time less than time t and the system completes its operation until time t starting from state S_1 . The third term means that the system moves from

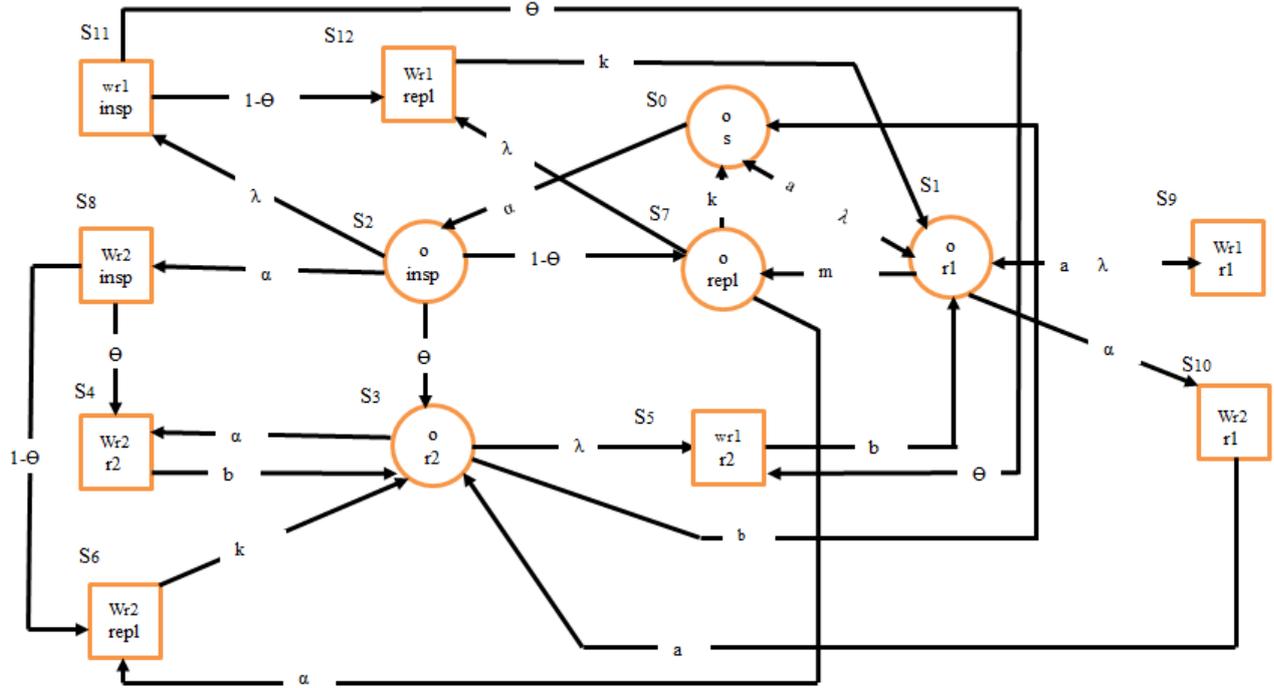


Fig. 1. Transition diagram ○ : Up state, □ : Down state.

state S_0 to state S_2 at time less than time t and the system completes its operation until time t starting from state S_2 . All other equations can be explained in the same manner.

$$\bar{\Pi}_1(t) = \bar{F}_1(t) \bar{F}_2(t) \bar{G}_1(t) \bar{M}(t) + q_{10}(t) \odot \bar{\Pi}_1(t) + q_{17}(t) \odot \bar{\Pi}_7(t), \quad (2)$$

$$\bar{\Pi}_2(t) = \bar{F}_1(t) \bar{F}_2(t) \bar{H}(t) + q_{23}(t) \odot \bar{\Pi}_3(t) + q_{27}(t) \odot \bar{\Pi}_7(t), \quad (3)$$

$$\bar{\Pi}_3(t) = \bar{F}_1(t) \bar{F}_2(t) \bar{G}_1(t) + q_{30}(t) \odot \bar{\Pi}_0(t), \quad (4)$$

$$\bar{\Pi}_7(t) = \bar{F}_1(t) \bar{F}_2(t) \bar{K}(t) + q_{70}(t) \odot \bar{\Pi}_0(t). \quad (5)$$

Making use of Laplace transform (LT) for Equations[1- 5] and solving for $\bar{\Pi}_0^*(u)$ considering $u = 0$, we have the time to system failures (MTSF) as follows:

$$MTSF = E(T) = \frac{l}{L}, \quad (6)$$

where,

$$l = \mu_0 + P_1(\mu_1 + P_{17}\mu_{17}) + P_2(\mu_2 + P_{23}\mu_3 + P_{27}\mu_7),$$

and

$$L = 1 - P_1(P_{10} + P_{17}P_{70}) - P_2(P_{23}P_{30} + P_{27}P_{70})$$

4. AVAILABILITY ANALYSES

Making use of probabilistic arguments, we obtain the following recursive relations for , $i = 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12$.

Where

$$AV_0(t) = M_0(t) + q_{01}(t) \odot AV_0(t) + q_{02}(t) \odot AV_2(t), \quad (7)$$

$$AV_1(t) = M_1(t) + q_{10}(t) \odot AV_0(t) + q_{17}(t) \odot AV_7(t) + q_{1,9}(t) \odot AV_9(t) + q_{1,10}(t) \odot AV_{10}(t), \quad (8)$$

$$AV_2(t) = M_2(t) + q_{23}(t) \odot AV_3(t) + q_{28}(t) \odot AV_8(t) + q_{27}(t) \odot AV_7(t) + q_{2,11}(t) \odot AV_{11}(t), \quad (9)$$

$$AV_3(t) = M_3(t) + q_{35}(t) \odot AV_5(t) + q_{30}(t) \odot AV_0(t) + q_{34}(t) \odot AV_4(t), \quad (10)$$

$$AV_4(t) = q_{43}(t) \odot AV_3(t), \quad (11)$$

$$AV_5(t) = q_{51}(t) \odot AV_1(t), \quad (12)$$

$$AV_6(t) = q_{63}(t) \odot AV_3(t), \quad (13)$$

$$AV_7(t) = M_7(t) + q_{70}(t) \odot AV_0(t) + q_{7,12}(t) \odot AV_{12}(t) + q_{76}(t) \odot AV_6(t), \quad (14)$$

$$AV_8(t) = q_{84}(t) \odot AV_4(t) + q_{86}(t) \odot AV_6(t), \quad (15)$$

$$AV_9(t) = q_{9,1}(t) \odot AV_1(t), \quad (16)$$

$$AV_{10}(t) = q_{10,3}(t) \odot AV_3(t), \quad (17)$$

$$AV_{11}(t) = q_{11,5}(t) \odot AV_5(t) + q_{11,12}(t) \odot AV_{12}(t), \quad (18)$$

$$AV_{12}(t) = q_{12,1}(t) \odot AV_1(t), \quad (19)$$

where, $M_0(t) = e^{-x(\alpha+\lambda)}$, $M_1(t) = e^{-x(\alpha+\lambda+a+m)}$, $M_2(t) = e^{-x(\alpha+\lambda+h)}$, $M_3(t) = e^{-x(\alpha+\lambda+b)}$, $M_7(t) = e^{-x(\alpha+\lambda+k)}$.

Taking LT of [7- 19] to solve them for $AV_0^*(u)$, then we get the steady state availability of the system AV_0 in the form ,

$$AV_0 = \lim_{u \rightarrow 0} u AV_0^*(u) = \frac{l_1}{L_1}, \quad (20)$$

5. BUSY PERIOD ANALYSIS

5.1 Expected Busy Period with Repair due to hardware failure:

Using probabilistic arguments we obtain

$$B_0^1(t) = q_{01}(t) \odot AV_0(t) + q_{02}(t) \odot AV_2(t), \quad (21)$$

$$B_1^1(t) = \bar{G}_1(t) + q_{10}(t) \odot AV_0(t) + q_{17}(t) \odot AV_7(t) + q_{1,10}(t) \odot AV_{10}(t) + q_{1,9}(t) \odot AV_9(t), \quad (22)$$

$$B_2^1(t) = q_{23}(t) \odot AV_3(t) + q_{28}(t) \odot AV_8(t) + q_{27}(t) \odot AV_7(t) + q_{2,11}(t) \odot AV_{11}(t), \quad (23)$$

$$B_3^1(t) = q_{35}(t) \odot AV_5(t) + q_{30}(t) \odot AV_0(t) + q_{34}(t) \odot AV_4(t), \quad (24)$$

$$B_4^1(t) = q_{43}(t) \odot AV_3(t), \quad (25)$$

$$B_5^1(t) = q_{51}(t) \odot AV_1(t), \quad (26)$$

$$B_6^1(t) = q_{63}(t) \odot AV_3(t), \quad (27)$$

$$B_7^1(t) = q_{70}(t) \odot AV_0(t) + q_{7,12}(t) \odot AV_{12}(t) + q_{76}(t) \odot AV_6(t), \quad (28)$$

$$B_8^1(t) = q_{84}(t) \odot AV_4(t) + q_{86}(t) \odot AV_6(t), \quad (29)$$

$$B_9^1(t) = \bar{G}_1(t) + q_{9,1}(t) \odot AV_1(t), \quad (30)$$

$$B_{10}^1(t) = \bar{G}_1(t) + q_{10,3}(t) \odot AV_3(t), \quad (31)$$

$$B_{11}^1(t) = q_{11,5}(t) \odot AV_5(t) + q_{11,12}(t) \odot AV_{12}(t), \quad (32)$$

$$B_{12}^1(t) = q_{12,1}(t) \odot AV_1(t), \quad (33)$$

Using LT to solve equations [21 – 33] for $B_0^*(u)$, we have the expected busy period with repair due to hardware failure in steady state as follows

$$B_0^1 = \lim_{u \rightarrow 0} u B_0^*(u) = \frac{l_2}{L_1}, \quad (34)$$

5.2 Expected Busy Period with Repair due to software failure:

Similarly as in (5.1), we can get the expected busy period with repair due to software failure in the steady state in the form,

$$B_0^2 = \lim_{u \rightarrow 0} u B_0^*(u) = \frac{l_3}{L_1}, \quad (35)$$

5.3 Expected Busy Period due to Replacment:

Similarly as in (5.1), we can get the expected busy period due to replacment in the steady state in the form,

$$B_0^3 = \lim_{u \rightarrow 0} u B_0^*(u) = \frac{l_4}{L_1}, \quad (36)$$

5.4 Expected Busy Period due to Inspection:

Similarly as in (5.1), we can get the expected busy period due to inspection in the steady state in the form,

$$B_0^4 = \lim_{u \rightarrow 0} u B_0^*(u) = \frac{l_5}{L_1}, \quad (37)$$

where

$$l_4(0) = \begin{pmatrix} 0 & A_{01} & A_{02} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & A_{17} & 0 & A_{19} & A_{110} & 0 & 0 \\ 0 & 0 & 1 & A_{23} & 0 & 0 & 0 & A_{27} & A_{28} & 0 & 0 & A_{211} & 0 \\ 0 & 0 & 0 & 1 & A_{34} & A_{35} & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ \frac{1}{k} & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ \frac{1}{k} & 0 & 0 & 0 & 0 & 0 & 0 & A_{76} & 1 & 0 & 0 & 0 & A_{712} \\ 0 & 0 & 0 & 0 & \theta & 0 & 1-\theta & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & \theta & 0 & 0 & 0 & 0 & 0 & 1 & 1-\theta \\ \frac{1}{k} & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix},$$

and

$$l_5(0) = \begin{pmatrix} 0 & A_{01} & A_{02} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & A_{17} & 0 & A_{19} & A_{110} & 0 & 0 \\ 0 & 0 & 1 & A_{23} & 0 & 0 & 0 & A_{27} & A_{28} & 0 & 0 & A_{211} & 0 \\ 0 & 0 & 0 & 1 & A_{34} & A_{35} & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ \frac{1}{h} & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & A_{76} & 1 & 0 & 0 & 0 & A_{712} \\ \frac{1}{h} & 0 & 0 & 0 & \theta & 0 & 1-\theta & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ \frac{1}{h} & 0 & 0 & 0 & 0 & \theta & 0 & 0 & 0 & 0 & 0 & 1 & 1-\theta \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}.$$

where,

$$\begin{aligned} A_{00} &= \frac{1}{\lambda+\alpha}, & A_{10} &= \frac{1}{\lambda+\alpha+a+m}, \\ A_{30} &= \frac{1}{\lambda+\alpha+b}, & A_{70} &= \frac{1}{\lambda+\alpha+k}, \\ A_{01} &= \frac{-\lambda}{\lambda+\alpha}, & A_{02} &= \frac{-\alpha}{\lambda+\alpha}, \\ A_{19} &= \frac{-\lambda}{\lambda+\alpha+a+m}, & A_{110} &= \frac{-\alpha}{\lambda+\alpha+a+m}, \\ A_{17} &= \frac{-m}{\lambda+\alpha+a+m}, & A_{211} &= \frac{-\lambda}{\lambda+\alpha+h}, \\ A_{28} &= \frac{-\alpha}{\lambda+\alpha+h}, & A_{27} &= (1-\theta) \frac{-h}{\lambda+\alpha+h}, \\ A_{23} &= (\theta) \frac{-h}{\lambda+\alpha+h}, & A_{35} &= \frac{-\lambda}{\lambda+\alpha+b}, \\ A_{34} &= \frac{-\alpha}{\lambda+\alpha+b}, & A_{712} &= \frac{-\lambda}{\lambda+\alpha+k}, \\ A_{76} &= \frac{-\alpha}{\lambda+\alpha+k}, \end{aligned}$$

and

$$\begin{aligned} B_{01} &= -q_{01}^*(u), & B_{02} &= -q_{02}^*(u), \\ B_{10} &= -q_{10}^*(u), & B_{17} &= -q_{17}^*(u), \\ B_{19} &= -q_{1,9}^*(u), & B_{110} &= -q_{1,10}^*(u), \\ B_{23} &= -q_{23}^*(u), & B_{27} &= -q_{27}^*(u), \\ B_{28} &= -q_{28}^*(u), & B_{211} &= -q_{2,11}^*(u), \\ B_{30} &= -q_{30}^*(u), & B_{34} &= -q_{34}^*(u), \\ B_{35} &= -q_{35}^*(u), & B_{43} &= -q_{43}^*(u), \\ B_{51} &= -q_{51}^*(u), & B_{63} &= -q_{63}^*(u), \\ B_{70} &= -q_{70}^*(u), & B_{76} &= -q_{76}^*(u), \\ B_{712} &= -q_{7,12}^*(u), & B_{84} &= -q_{84}^*(u), \\ B_{86} &= -q_{86}^*(u), & B_{91} &= -q_{9,1}^*(u), \\ B_{103} &= -q_{10,3}^*(u), & B_{115} &= -q_{11,5}^*(u), \\ B_{1112} &= -q_{11,12}^*(u), & B_{121} &= -q_{12,1}^*(u). \end{aligned}$$

6. COST BENEFIT ANALYSES

Since

$$C(t) = \varepsilon \mu_{up}(t) - \sum_{i=1}^4 K_i \mu_R^i(t) \quad (38)$$

where,

ε : revenue per unit of up time

K_1, K_2, K_3, K_4 : cost per unit time of repair due to hardware, software failures, replacement and inspection respectively.

$$\mu_{up}(t) = \int_0^t AV_0(u) du \quad (39)$$

$$\mu_R^i = \int_0^t B_0^i(u) du \quad (40)$$

Using (38), (39) and (40), we have

$$C^*(t) = \varepsilon \mu_{up}^*(t) - \sum_{i=1}^4 K_i \mu_R^{i*}(t) \quad (41)$$

Since the expected profit per unit of time in steady state is,

$$C = \lim_{t \rightarrow \infty} \frac{C(t)}{t} = \lim_{u \rightarrow 0} u^2 C^*(u) \quad (42)$$

Then from (42), we get

$$C = (\varepsilon l_1 - K_1 l_2 - K_2 l_3 - K_3 l_4 - K_4 l_5) / L_1 \quad (43)$$

where

ε : is the revenue per unit up-time of the system,

K_1 : is the cost per unit time for which the server is busy due to hardware failure,

K_2 : is the cost per unit time for which the server is busy due to software failure,

K_3 : is the cost per unit time for which the server is busy due to replacement,

K_4 : is the cost per unit time for which the server is busy due to inspection.

7. NUMERICAL EXAMPLE

By setting $\varepsilon = 2000$, $K_1 = 30$, $K_2 = 40$, $K_3 = 100$, $K_4 = 5$, figures display the variation of AV, MTSF, and profit respectively, for different values of a , b , k , m , h , α and λ .

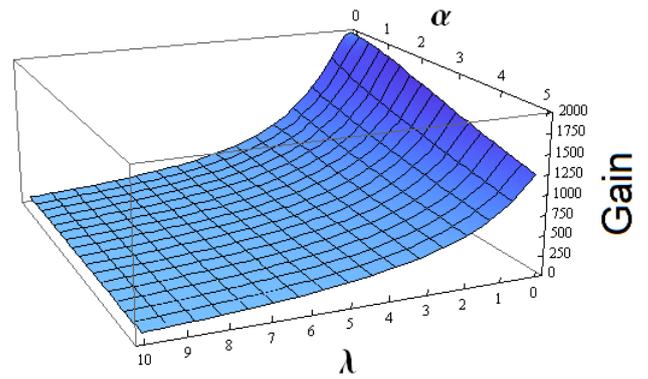


Fig. 4. Variation of Profit Gain (GT) with α, λ , ($a = 2.5, b = 5, k = 1.2, m = 30, h = 40, \theta = 0.98$)

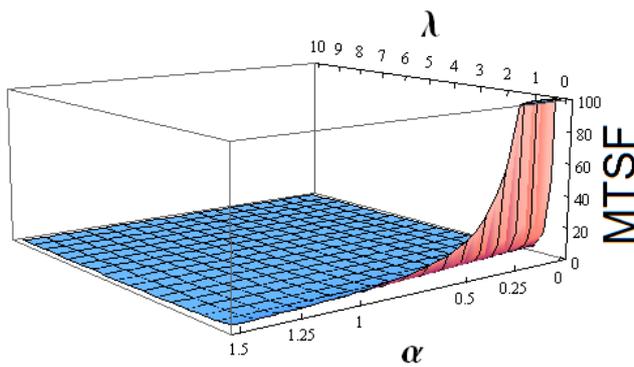


Fig. 2. Variation of Mean Time to System failure (MTSF) with α, λ ($a = 2.5, b = 5, k = 1.2, m = 30, h = 40, \theta = 0.98$)

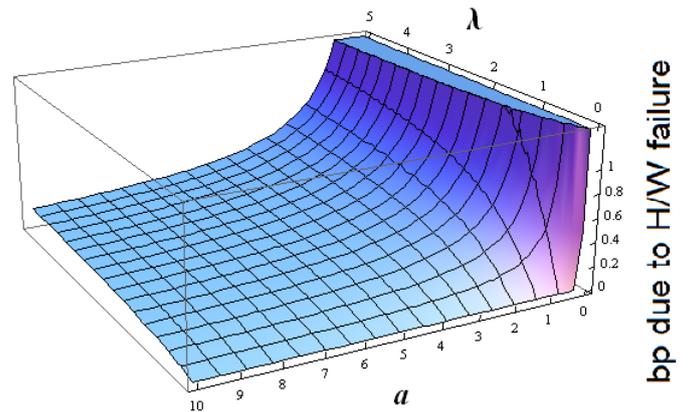


Fig. 5. Variation of Expected Busy Period due to hardware failure (bp) with λ, a , ($b = 5, k = 1.2, m = 30, h = 40, \alpha = 0.02, \theta = 0.98$)

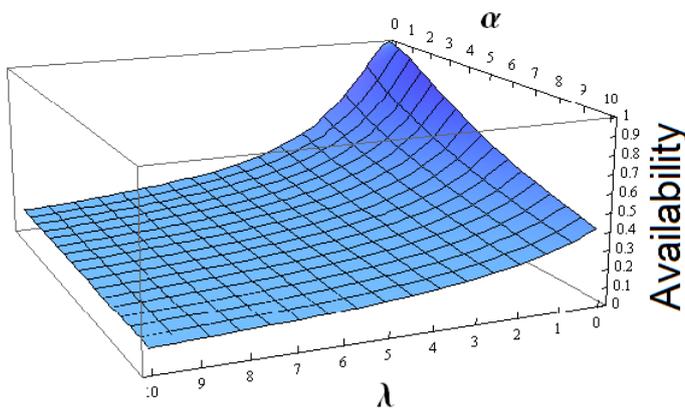


Fig. 3. Variation of Availability (AV) with α, λ ($a = 2.5, b = 5, k = 1.2, m = 30, h = 40, \theta = 0.98$)

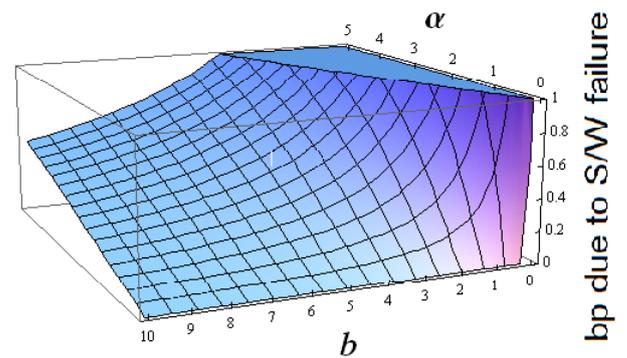


Fig. 6. Variation of Expected Busy Period due to software failure (bp) with α, b , ($a = 2.5, k = 1.2, m = 30, h = 40, \lambda = 0.2, \theta = 0.98$)

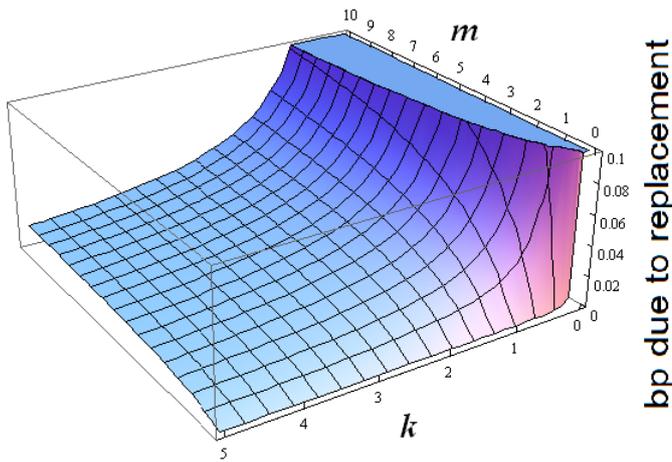


Fig. 7. Variation of Expected Busy Period due to replacement (bp) with m, k , ($a = 2.5, b = 5, h = 40, \lambda = 0.1, \alpha = 0.02, \theta = 0.98$)

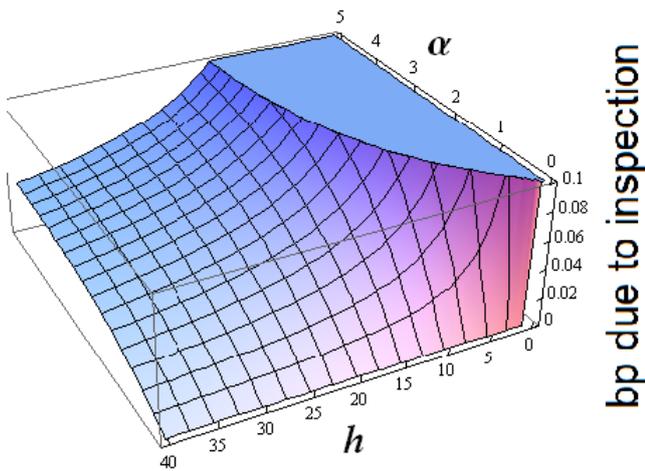


Fig. 8. Variation of Expected Busy Period due to Inspection (bp) with α, h , ($a = 2.5, b = 5, k = 1.2, m = 30, \lambda = 0.1, \theta = 0.98$)

8. CONCLUSION

In this paper we successfully obtained some reliability measures of the system such as, the MTSF, the steady state availability, the expected busy periods of the serverman and the expected profit of the system considering hardware failure, software failure, replacement and inspection using regenerative point technique .

1. The MTSF, AV and GT increase with decreasing of λ .
2. The MTSF, AV and GT increase with decreasing of α .
3. The busy period due to h/w failure increase with increasing λ and decreasing with increasing of a .
4. The busy period due to s/w failure increase with increasing α and decreasing with increasing of b .
5. The busy period due to replacement increase with increasing m and decreasing with increasing of k .
6. The busy period due to inspection increase with increasing α and decreasing with increasing of h .

According to the results obtained for a particular case, it is concluded that the system model can be more reliable and profitable to

use by increasing the repair and inspection rates of the unit and by making immediate replacement of the unit after the completion of maximum repair time.

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